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Combinatorial optimization with Rydberg platforms: advances and challenges

Colloque Atomes de Rydberg et simulation quantique
Friday, April 5th, 2024

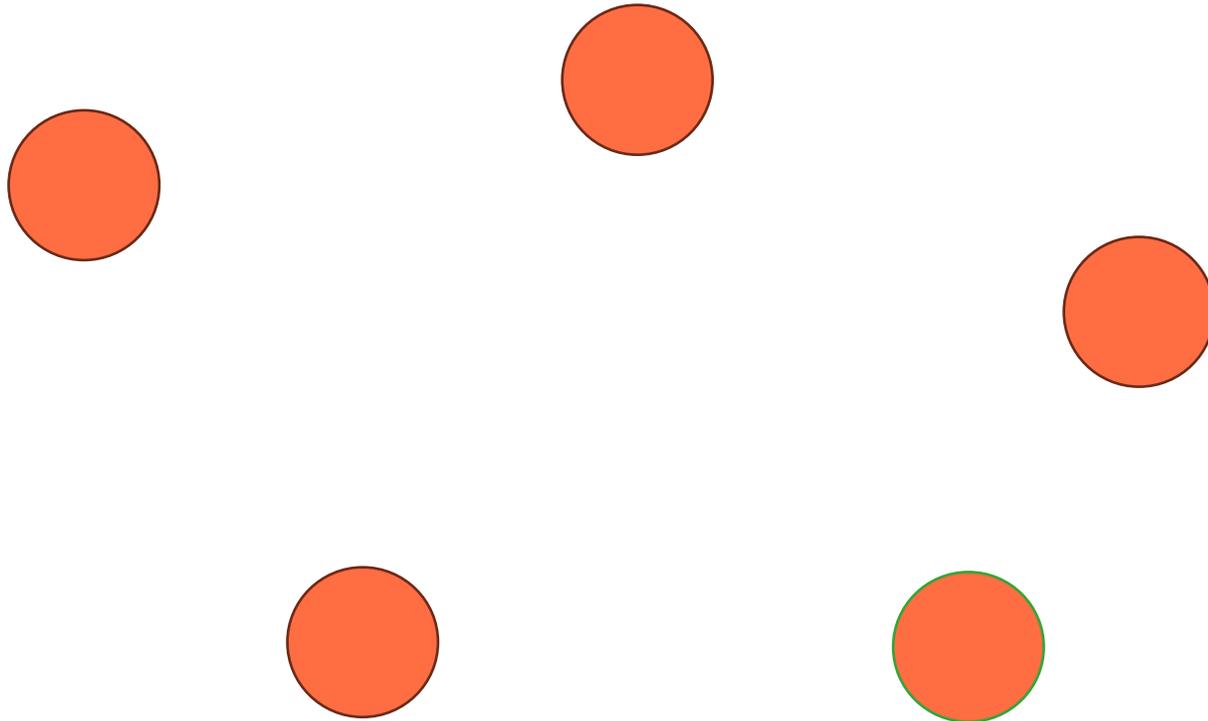
Thomas Ayrat
Eviden Quantum Lab, France

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an atos business

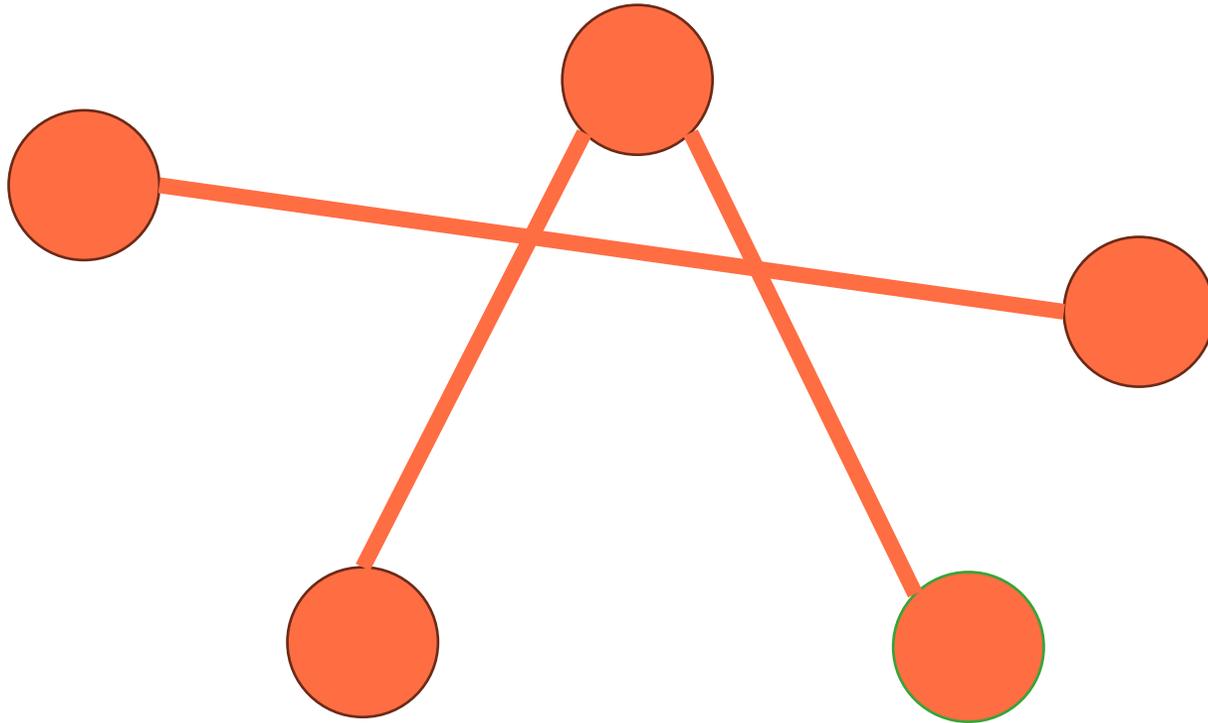
An example of combinatorial optimization problem

A successful party



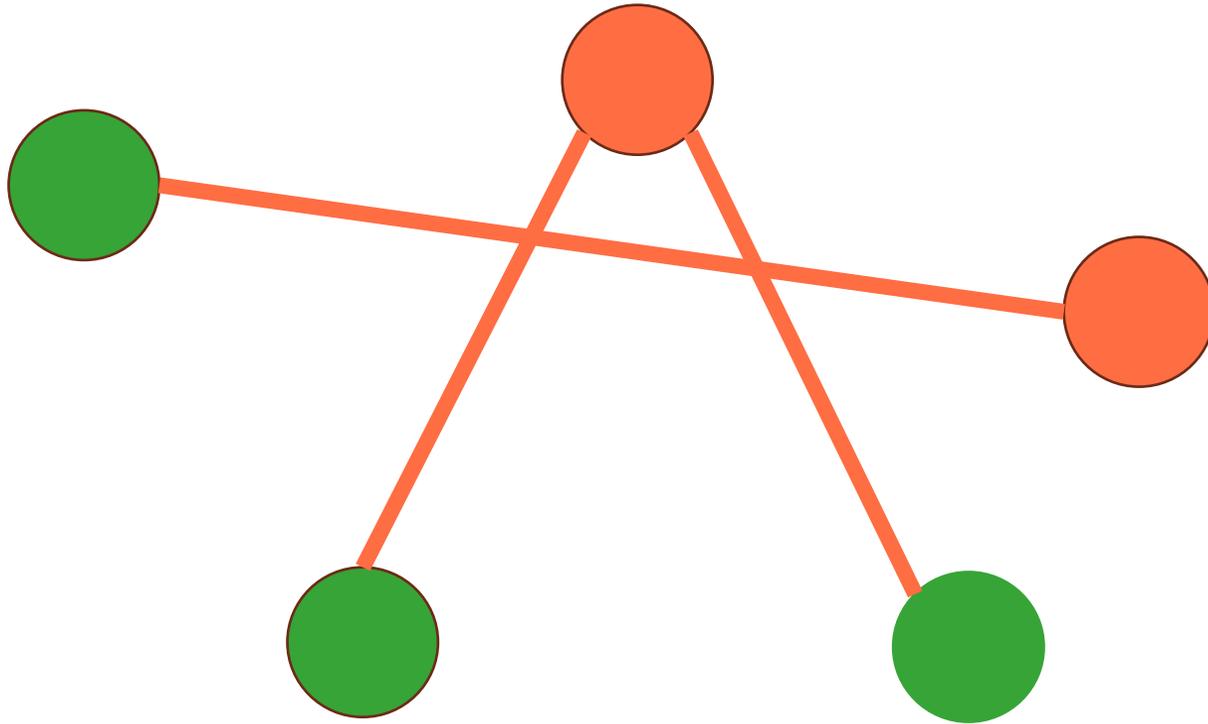
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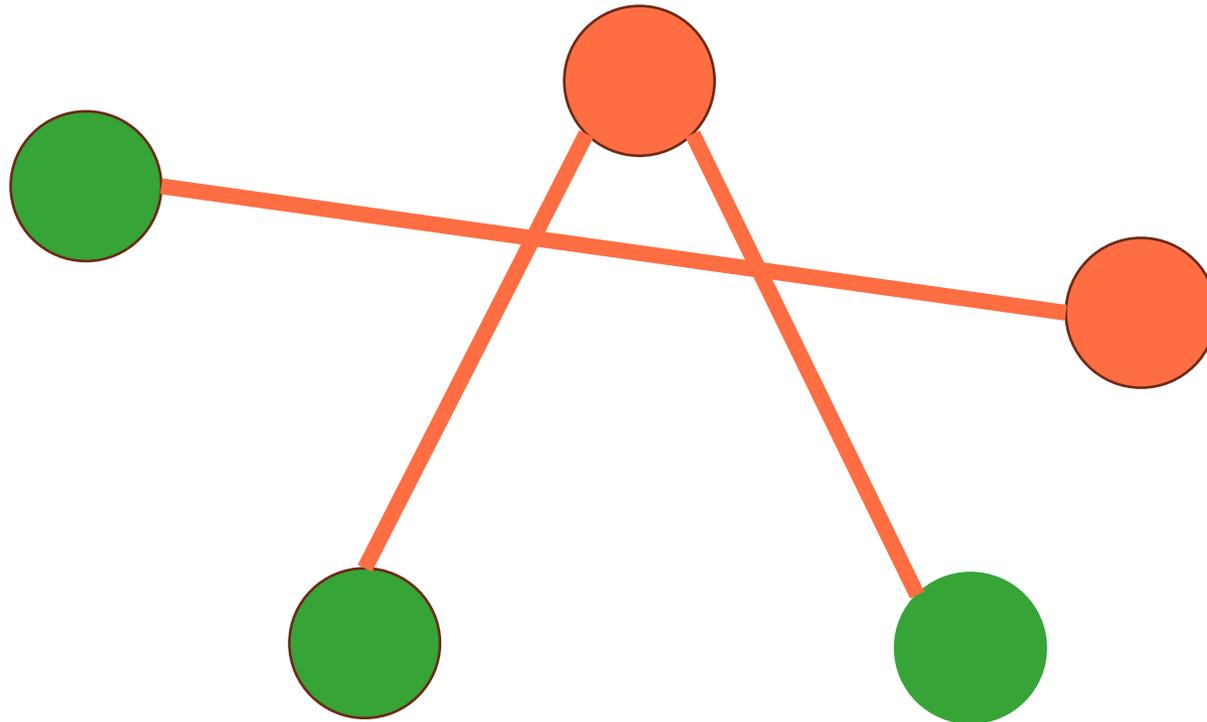
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A successful party

**Maximum independent set (MIS)
problem**



Optimal solution?

runtime exponential with system size!

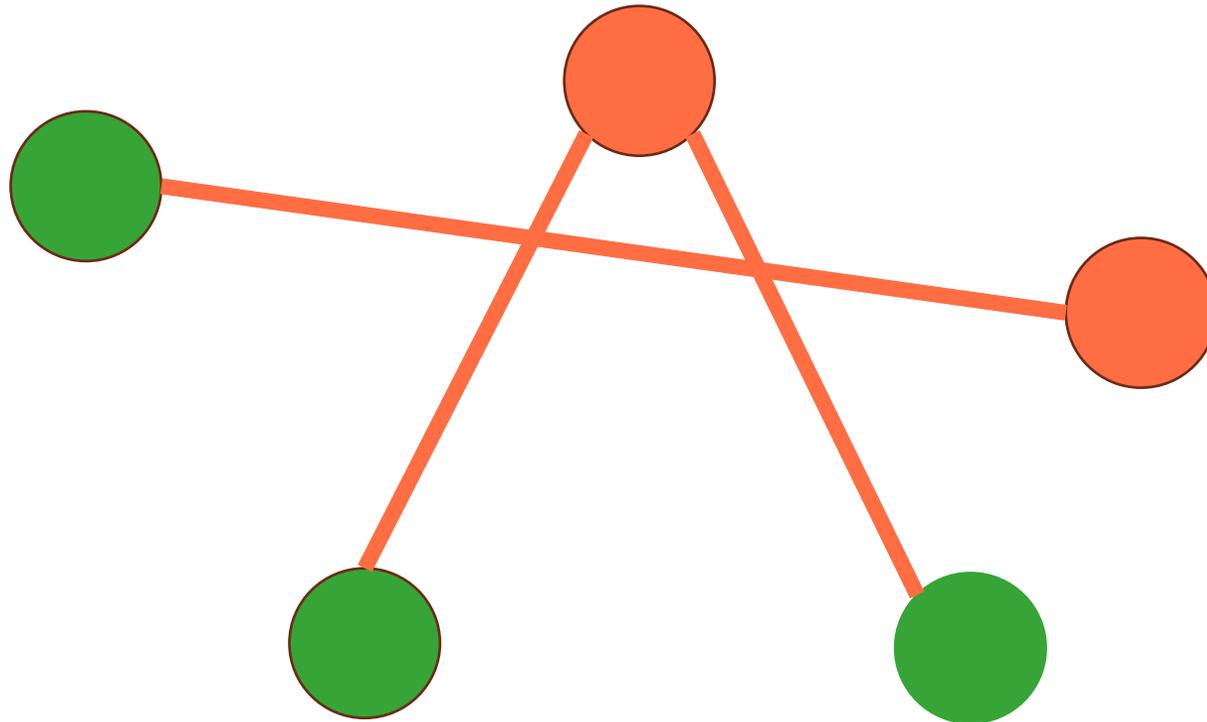
Approximate solution?

hard limit to get close to optimal solution!

An example of combinatorial optimization problem

A successful party

Maximum independent set (MIS) problem



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runtime exponential with system size!

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hard limit to get close to optimal solution!

Many other examples

Traveling salesperson, maximum cut problem...

Many industrial applications!

Can (Rydberg) quantum processors help?

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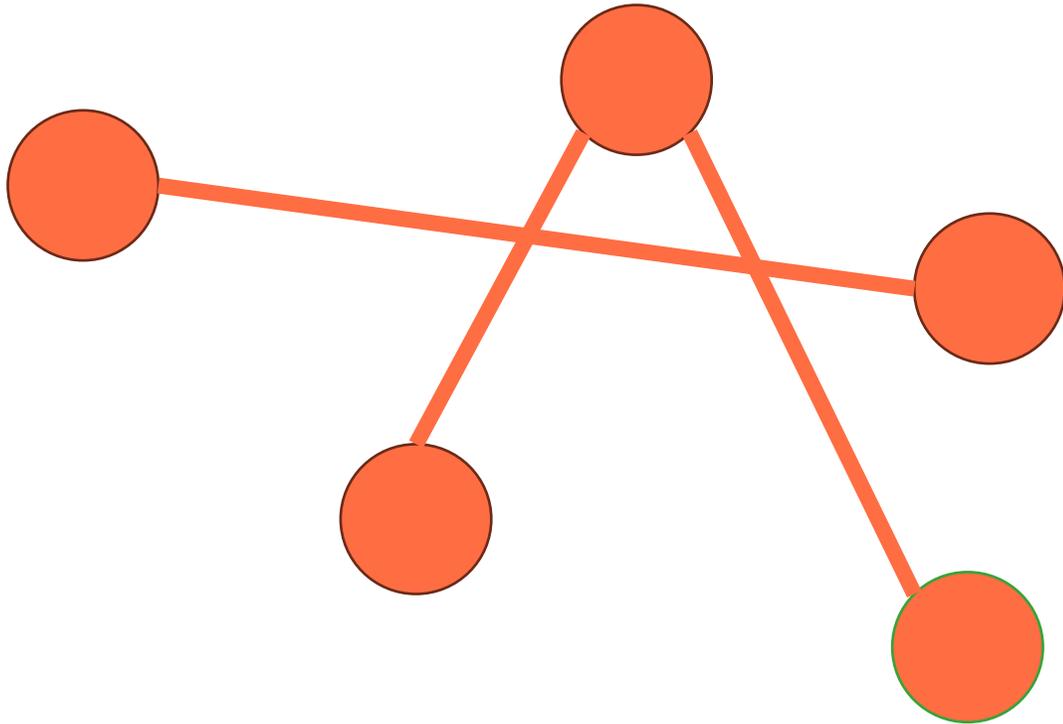
Outline

- 1 The unit-disk maximum independent set problem: a promising application for Rydberg quantum simulators?
- 2 The challenge of decoherence
- 3 Towards more general graphs... and applications

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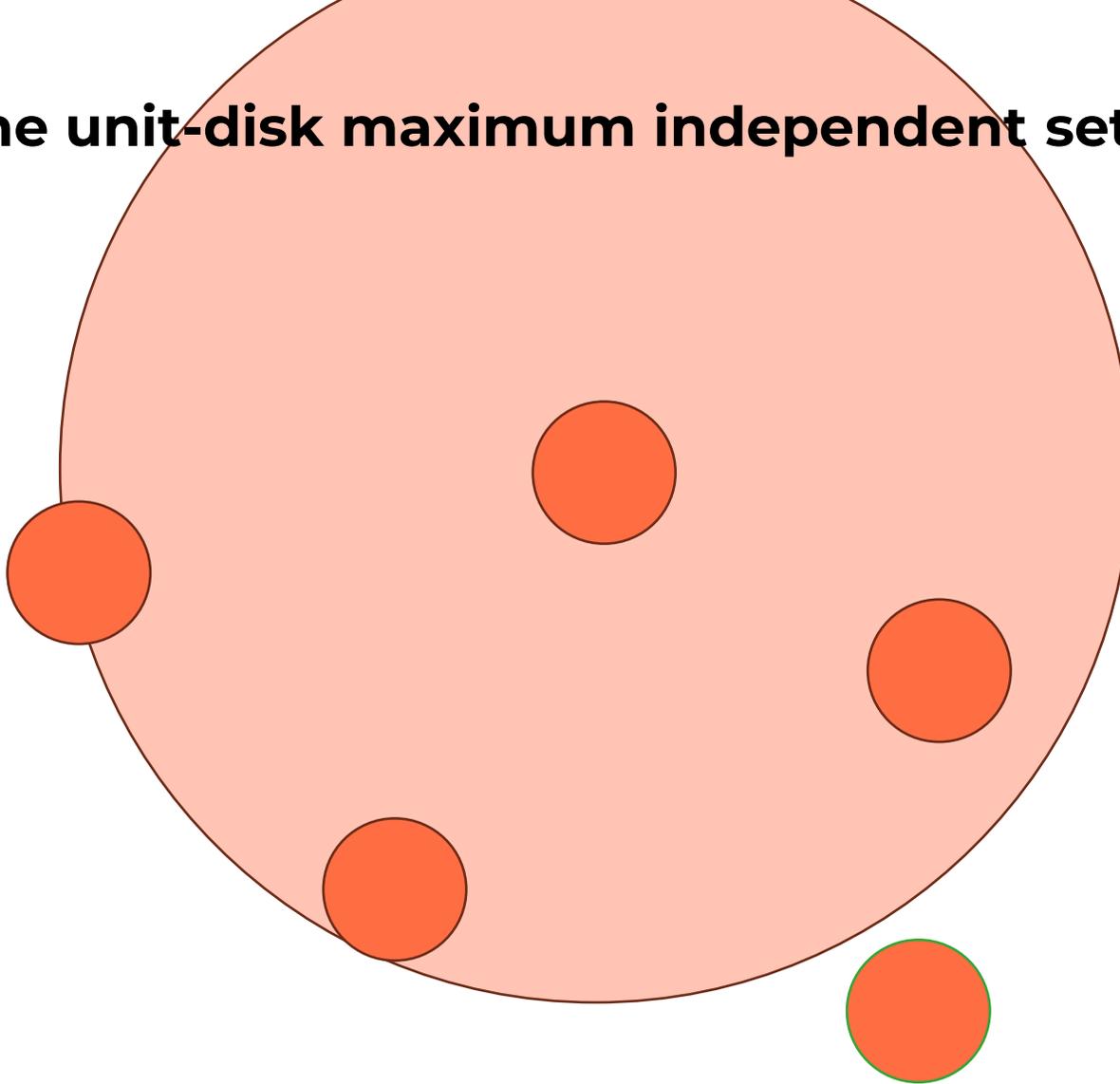
- 1 The unit-disk maximum independent set problem: a promising application for Rydberg quantum simulators?

The unit-disk maximum independent set



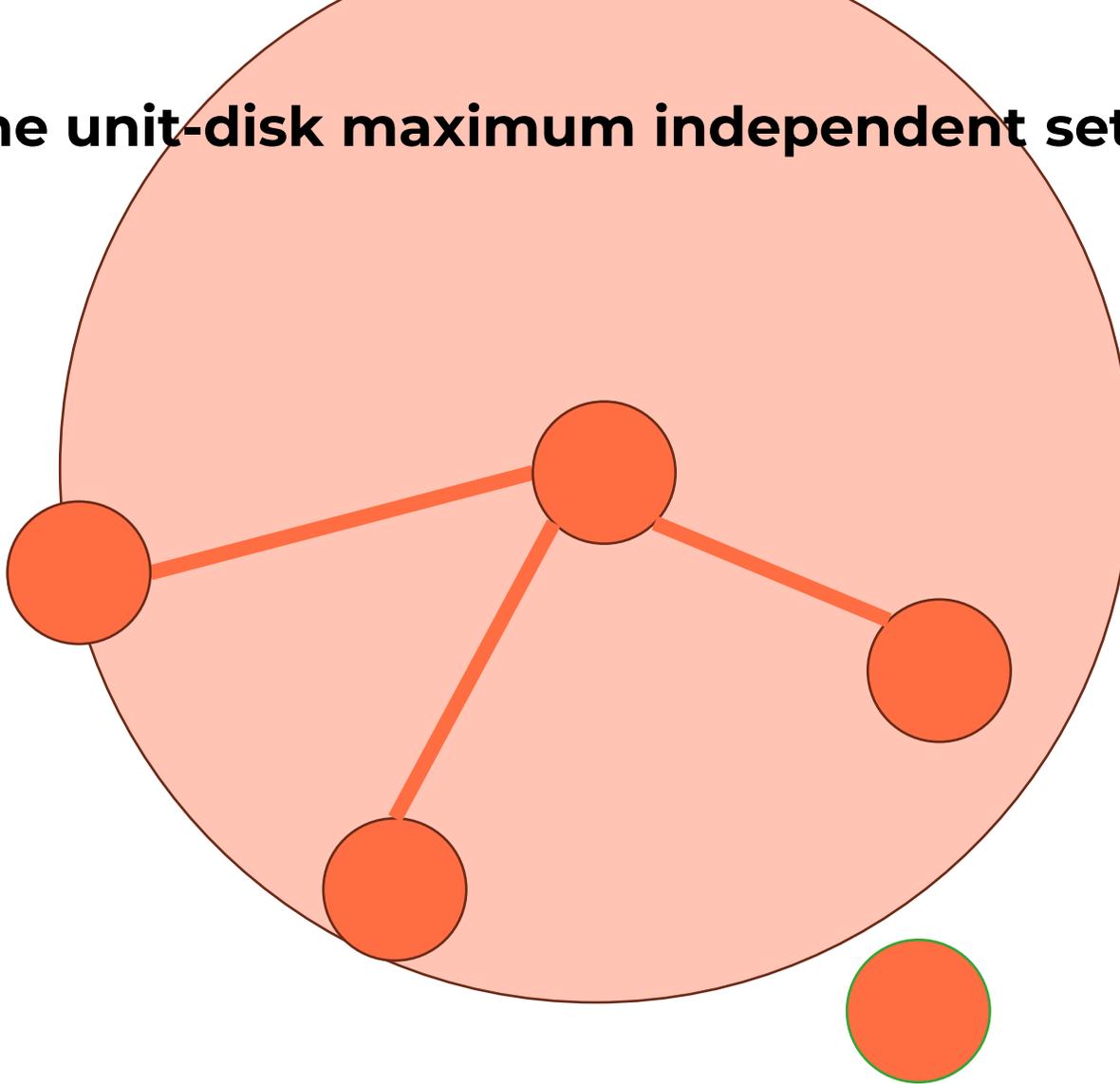
The unit-disk maximum independent set

Restriction to unit-disk graphs: **UDMIS**



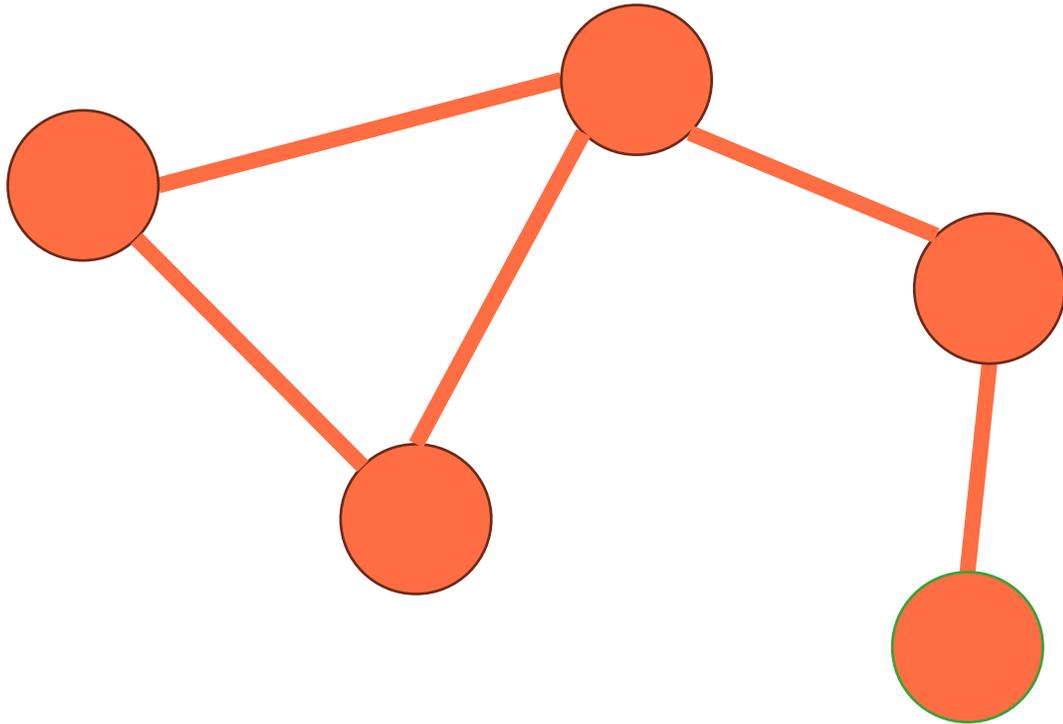
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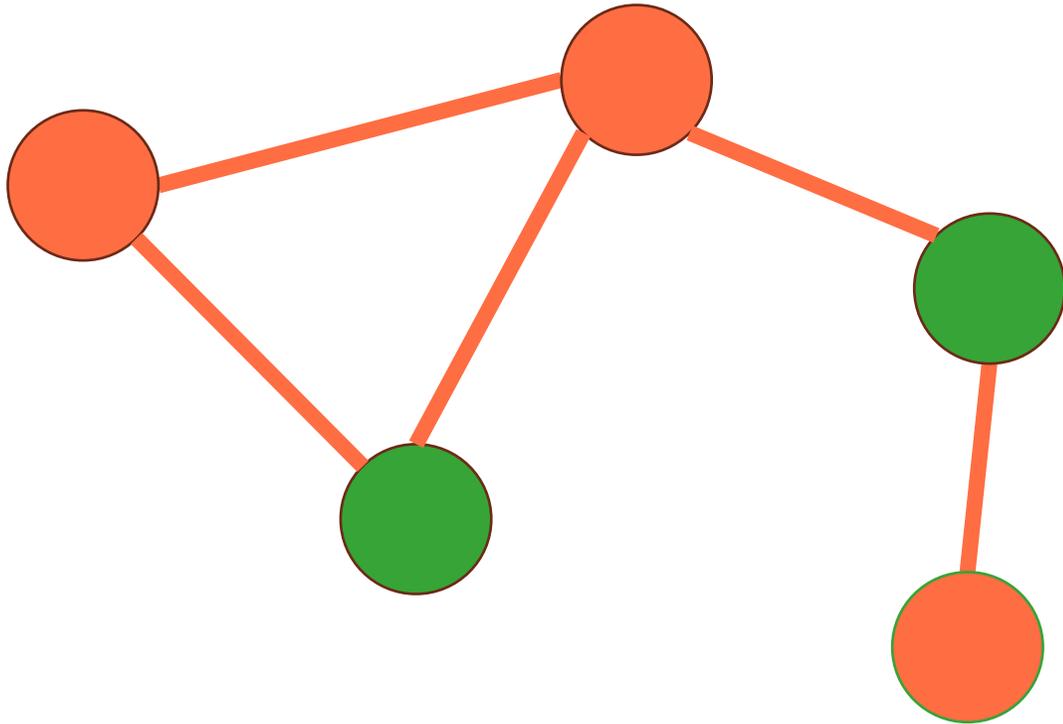


The unit-disk maximum independent set

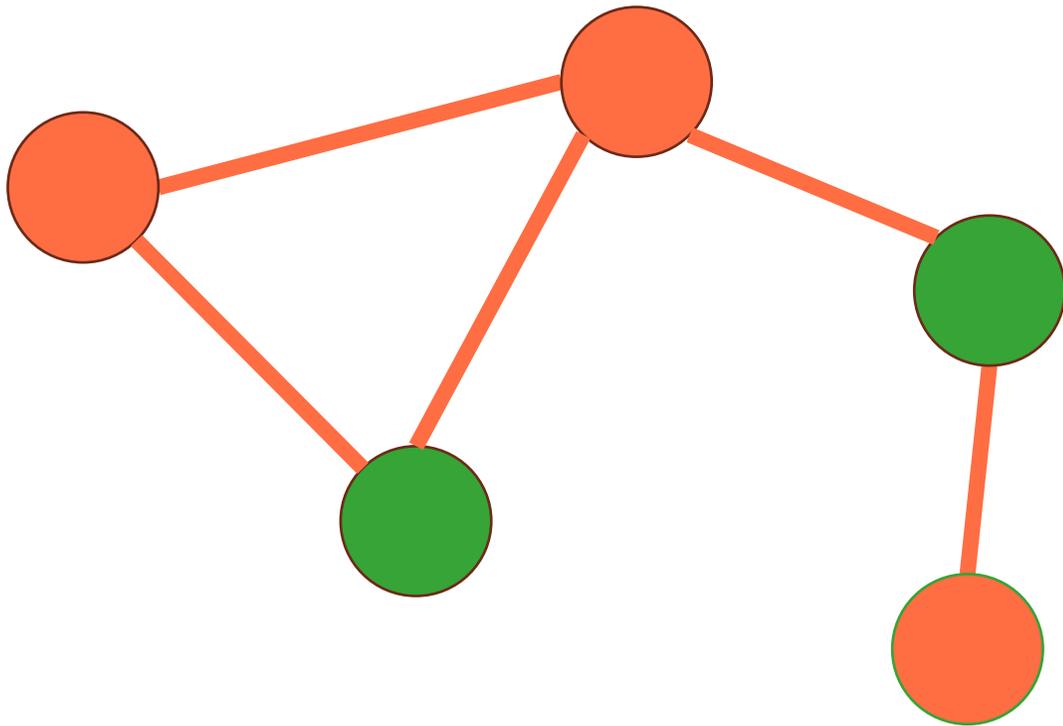
Restriction to unit-disk graphs: **UDMIS**

Still an exponential runtime for optimal solution.

For approximate solution:
can get ϵ close to optimal, but **runtime is exponential in $1/\epsilon$!**



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Precise definition of success?

Approximation ratio:

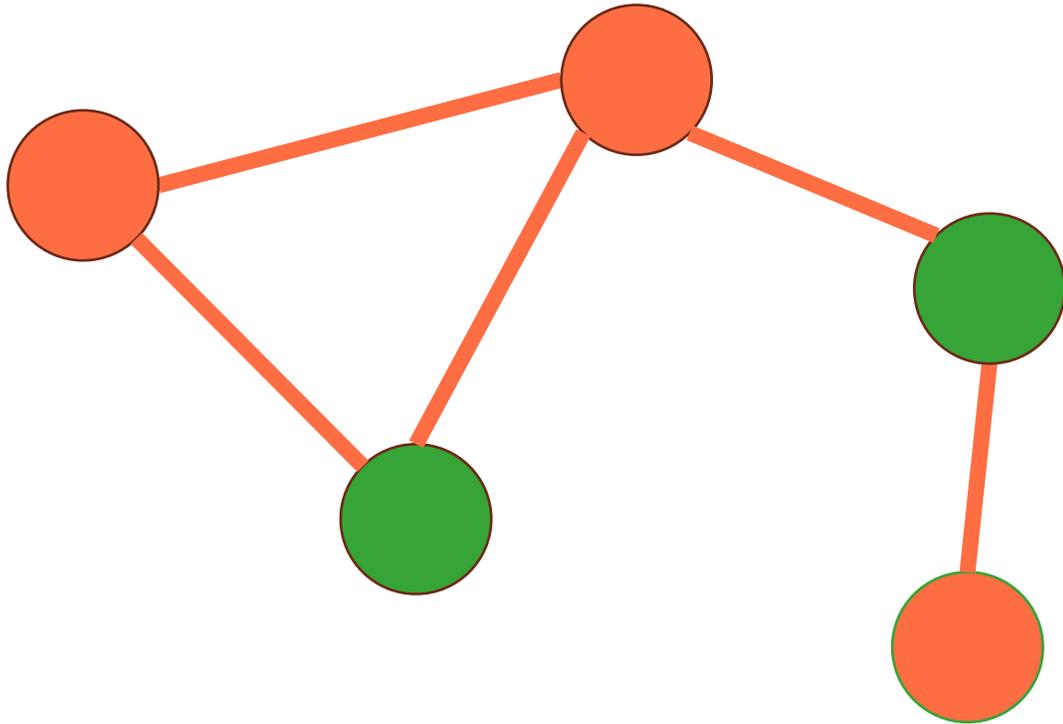
Your solution

Cost function \rightarrow

$$\alpha = \frac{C(S)}{C(S^*)} \leq 1$$

Optimal solution

The unit-disk maximum independent set



Restriction to unit-disk graphs: **UDMIS**

Still an exponential runtime for optimal solution.

For approximate solution:
can get ϵ close to optimal, but **runtime is exponential in $1/\epsilon$!**

$$\alpha = 1 - \epsilon$$

Precise definition of success?

Approximation ratio:

Your solution

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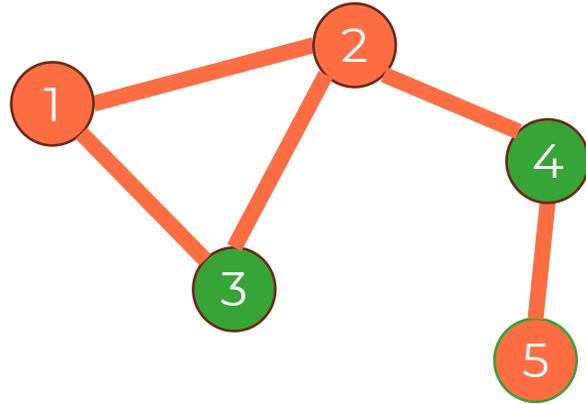
Optimal solution

How could quantum mechanics help?

Turn cost function minimization into estimation of ground state of Hamiltonian!

One solution = a string of bits $S = (n_1, n_2, \dots, n_n)$

Here $S = (0,0,1,1,0)$

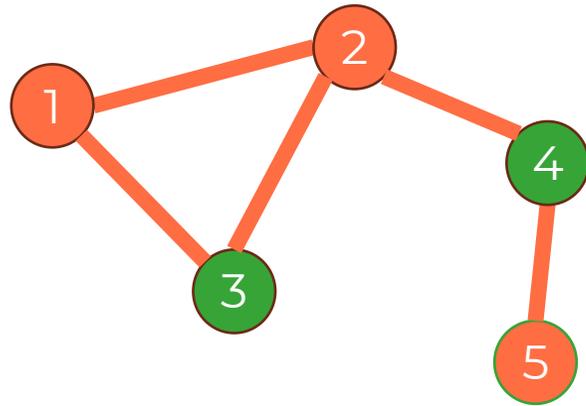


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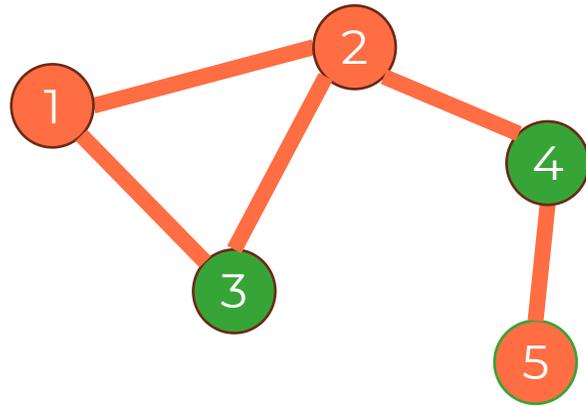
with constraint: $n_i n_j = 0$ if (i, j) is an edge

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To convert to Hamiltonian: easier to relax constraint with Lagrange multiplier:

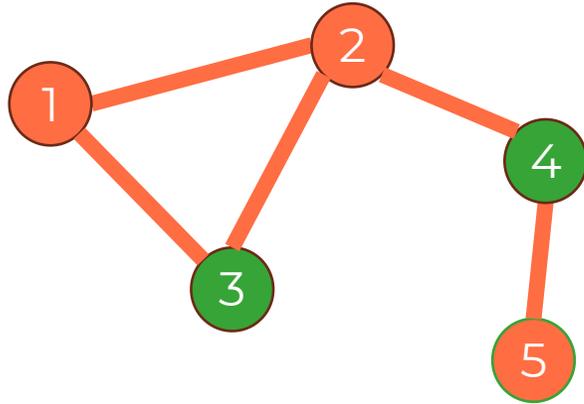
$$C(S, U) = \sum_i n_i - U' \sum_{i,j \in E} n_i n_j$$

How could quantum mechanics help?

Turn cost function minimization into estimation of ground state of Hamiltonian!

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To convert to Hamiltonian: easier to relax constraint with Lagrange multiplier:

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Turn into operators (+minus sign):

$$H = -\delta \sum_i \hat{n}_i + U \sum_{i,j \in E} \hat{n}_i \hat{n}_j$$

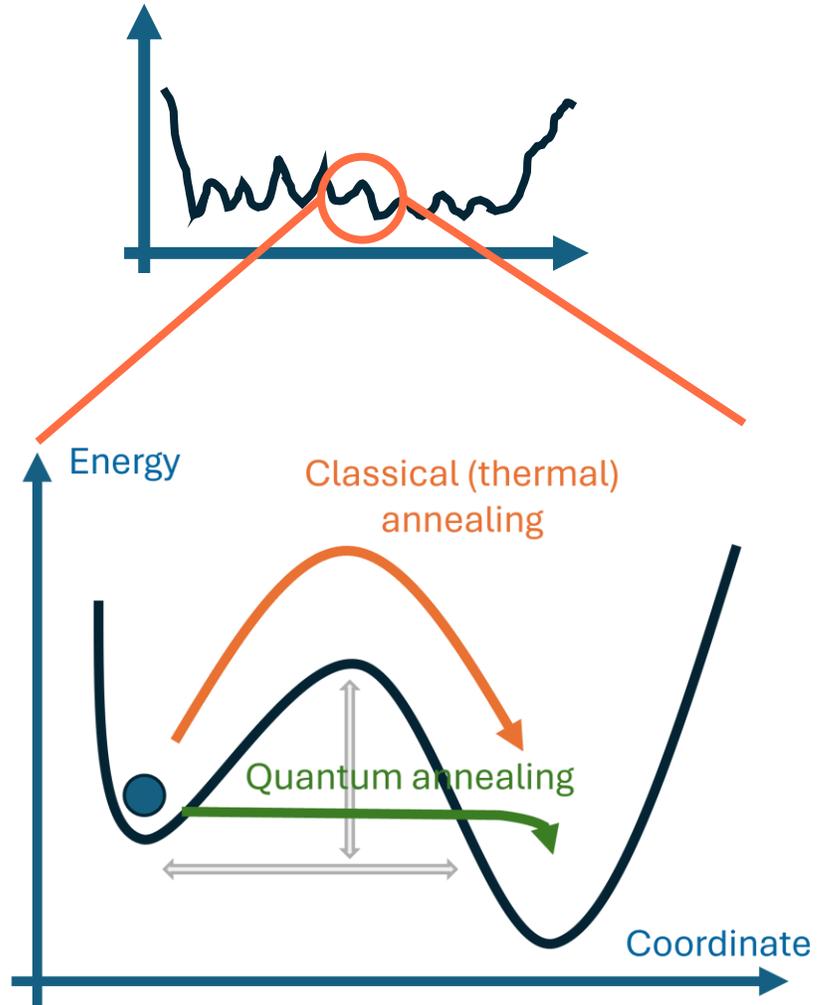
with $\hat{n}|0\rangle = 0, \hat{n}|1\rangle = |1\rangle$.

Equivalent quantum problem:

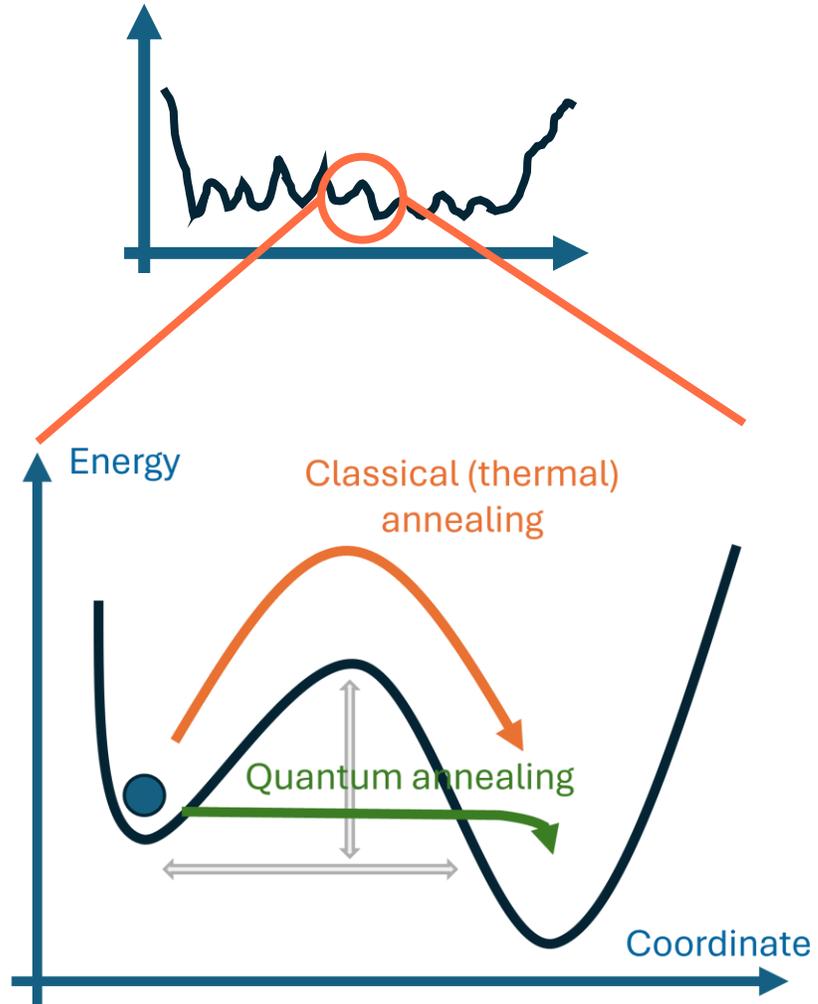
Find lowest eigenstate

$$H|\Psi_0\rangle = E_0|\Psi_0\rangle$$

Finding a quantum ground state in practice: quantum annealing



Finding a quantum ground state in practice: quantum annealing



To create quantum tunneling:

$$H(t) = \frac{t}{t_f} H + \left(1 - \frac{t}{t_f}\right) H_{\text{tunnel}}$$

with e.g

$$H_{\text{tunnel}} = \Omega \sum_i \hat{\sigma}_i^x$$

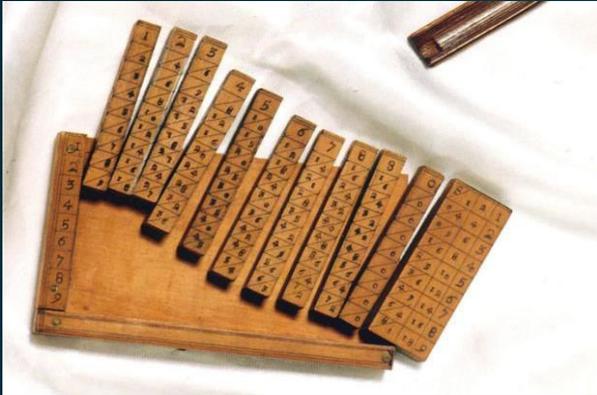
If annealing time t_f long enough (adiabatic):

Start from GS $|\Phi_0\rangle$ of H_{tunnel} , end in GS $|\Psi_0\rangle$ of H .

One experimental implementation: **d-wave computers: Classical + quantum annealing.**

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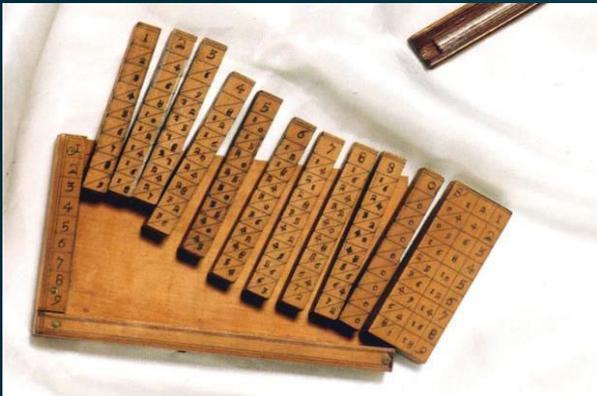
Analog computers?



Slide rule
(J. Napier, 1614)

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Analog computers?



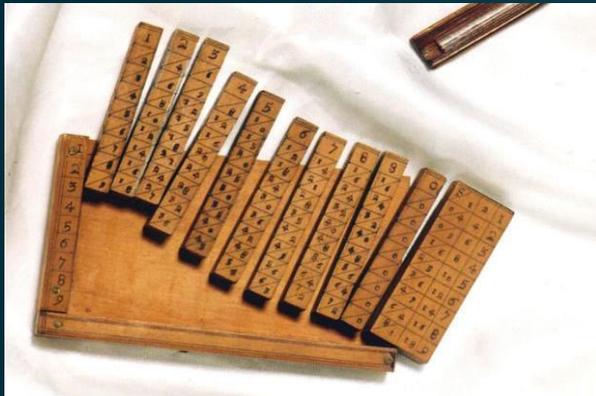
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Pascaline
(B. Pascal, 1642)

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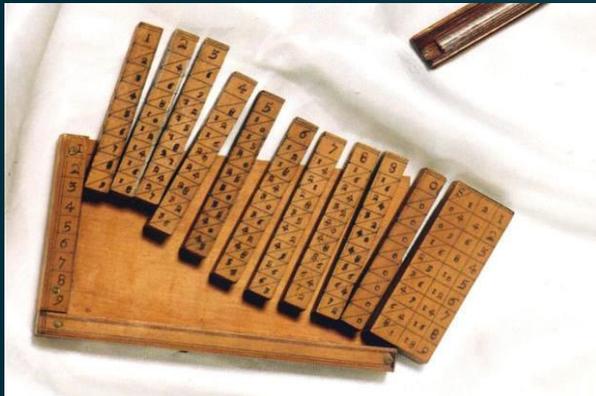
Bull machine
(F. Bull, 1920s)



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Analog computers?



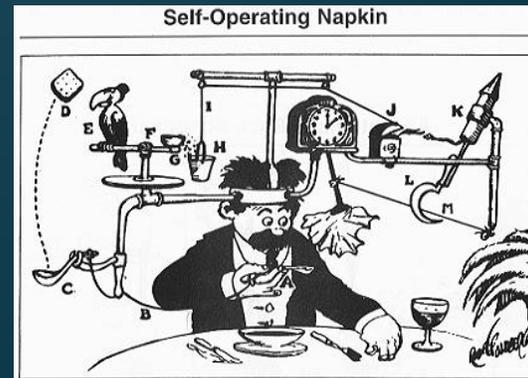
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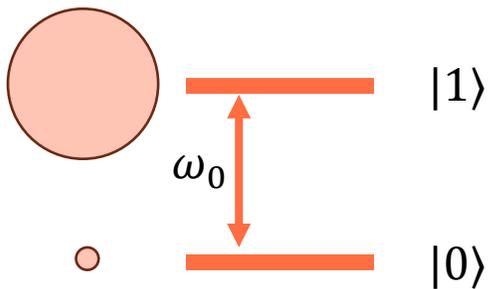
Rube Goldberg
machine (!)

Rydberg atoms: an analog quantum computer

We want to realize

$$H(t) = \frac{t}{t_f} \left(-\delta \sum_i \hat{n}_i + U \sum_{i,j \in E} \hat{n}_i \hat{n}_j \right) + \left(1 - \frac{t}{t_f} \right) \left(\Omega \sum_i \hat{\sigma}_i^x \right)$$

Rydberg atoms: artificial system that realizes a similar Hamiltonian

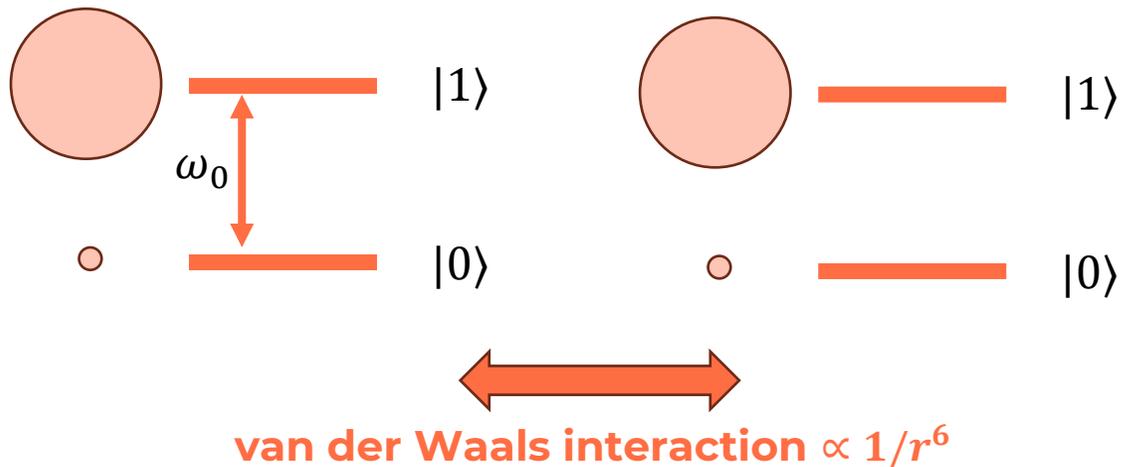


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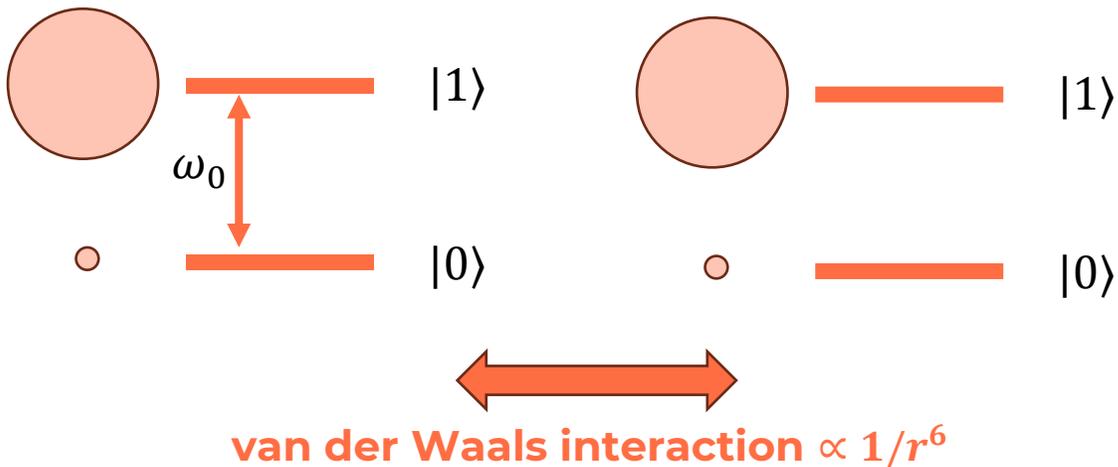
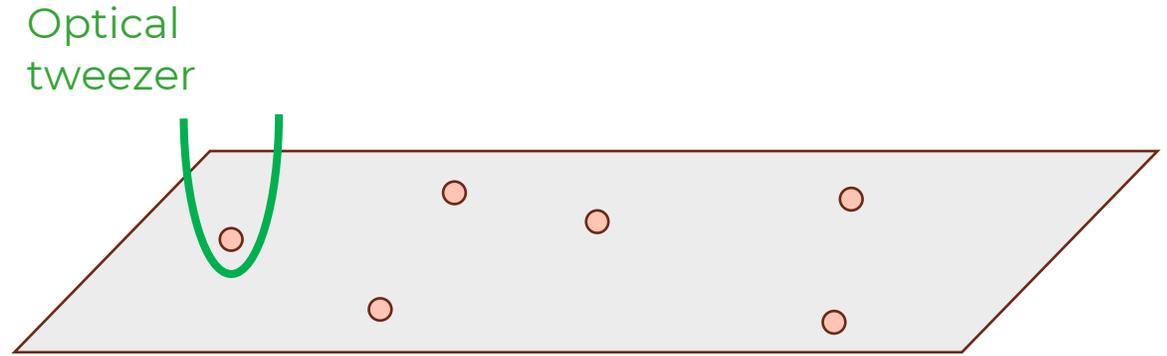


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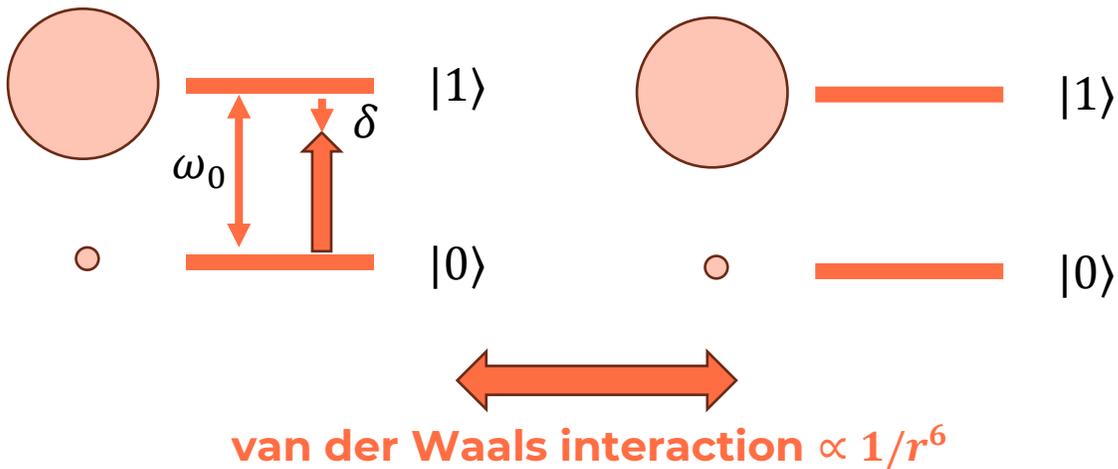
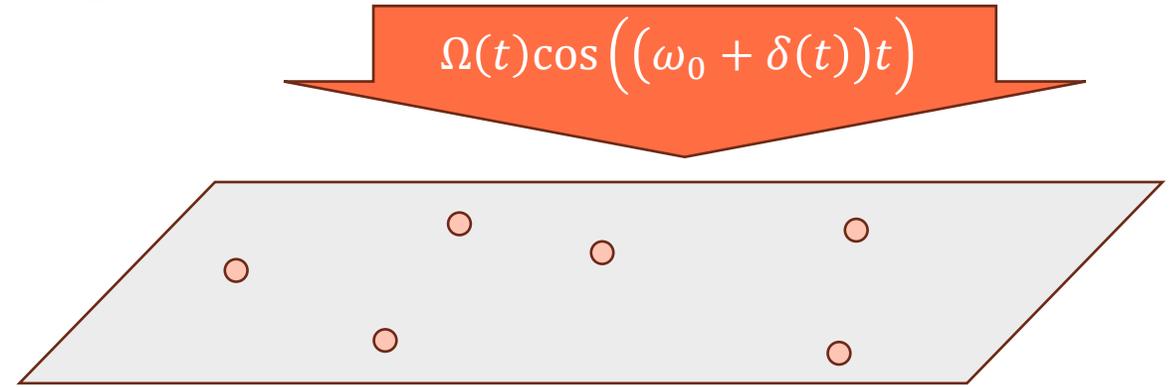


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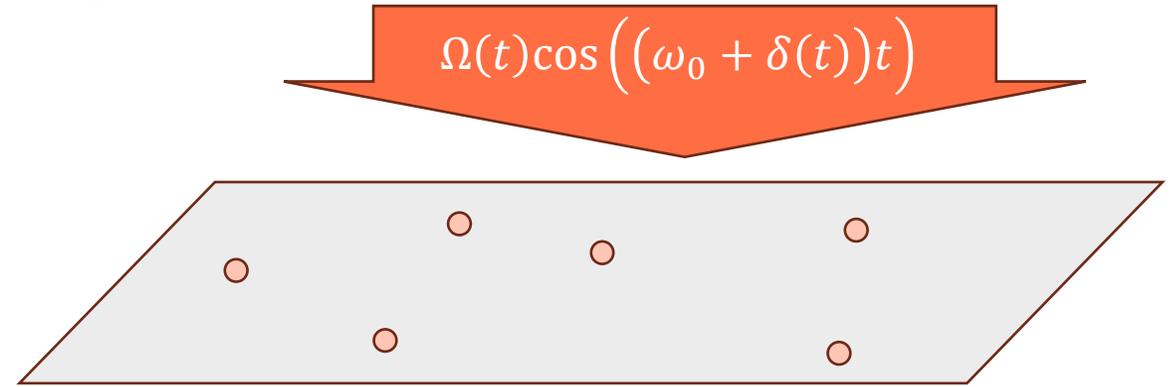
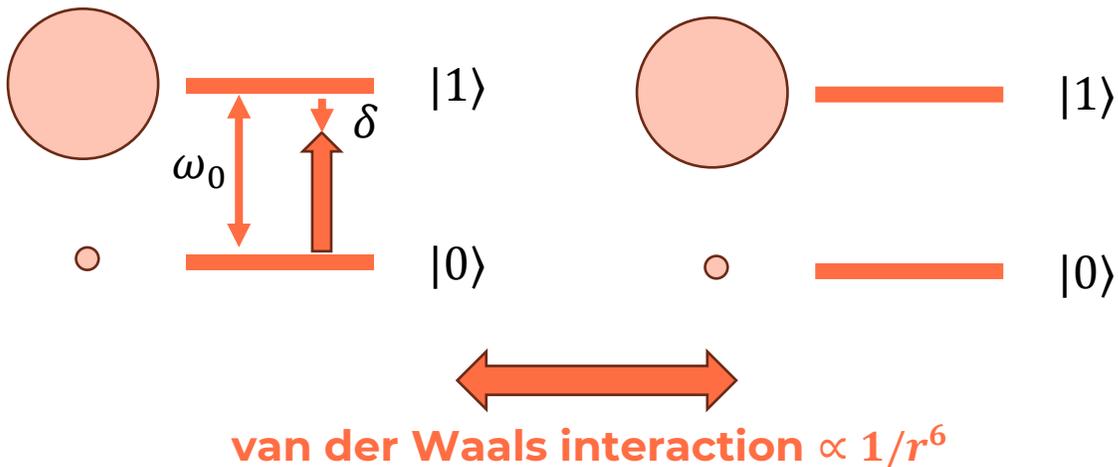


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Rydberg Hamiltonian:

$$H_{\text{Ryd}} = \underbrace{-\delta(t)}_{\text{detuning}} \sum_i \hat{n}_i + \sum_{i < j} \underbrace{\frac{C}{(r_i - r_j)^6}}_{\text{vdW}} \hat{n}_i \hat{n}_j + \frac{\Omega(t)}{2} \sum_i \hat{\sigma}_i^x$$

detuning

vdW

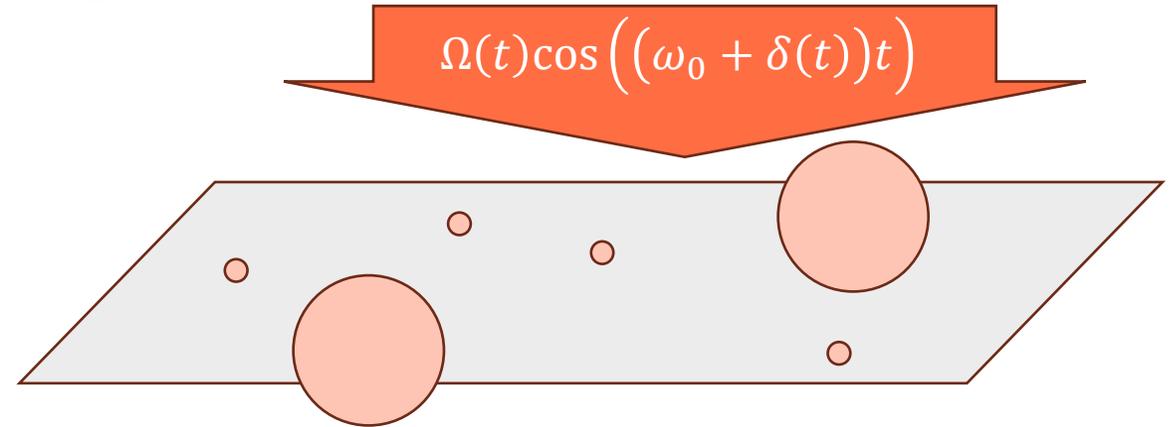
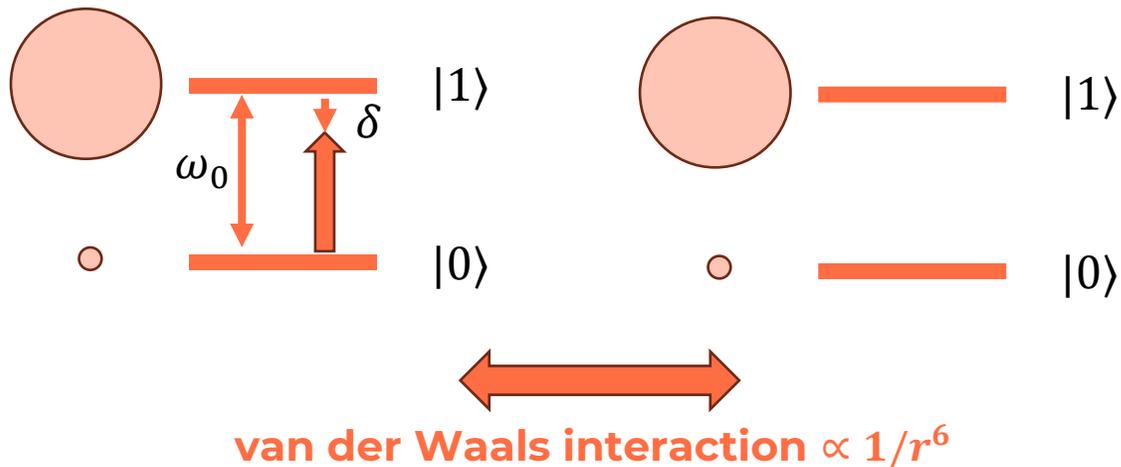
Rabi

Rydberg atoms: an analog quantum computer

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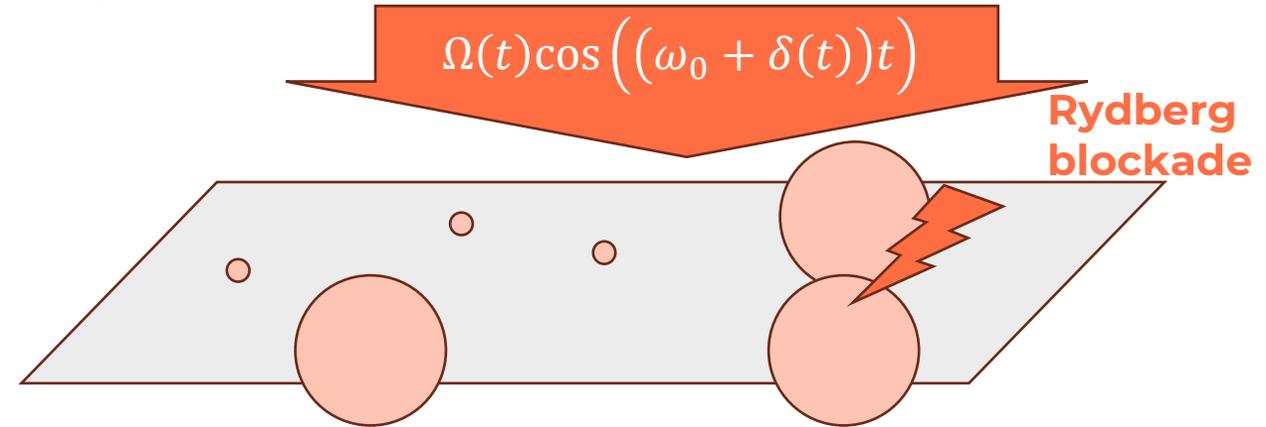
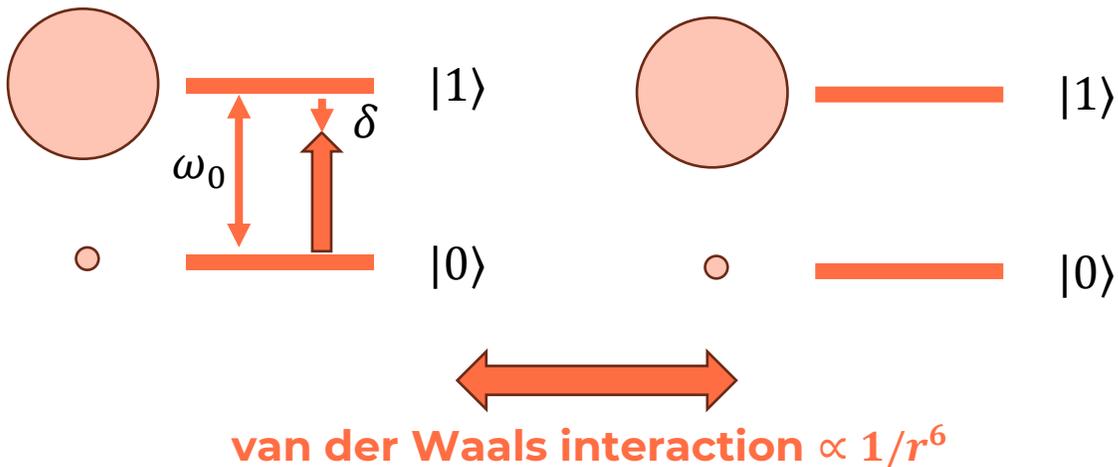
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vdW

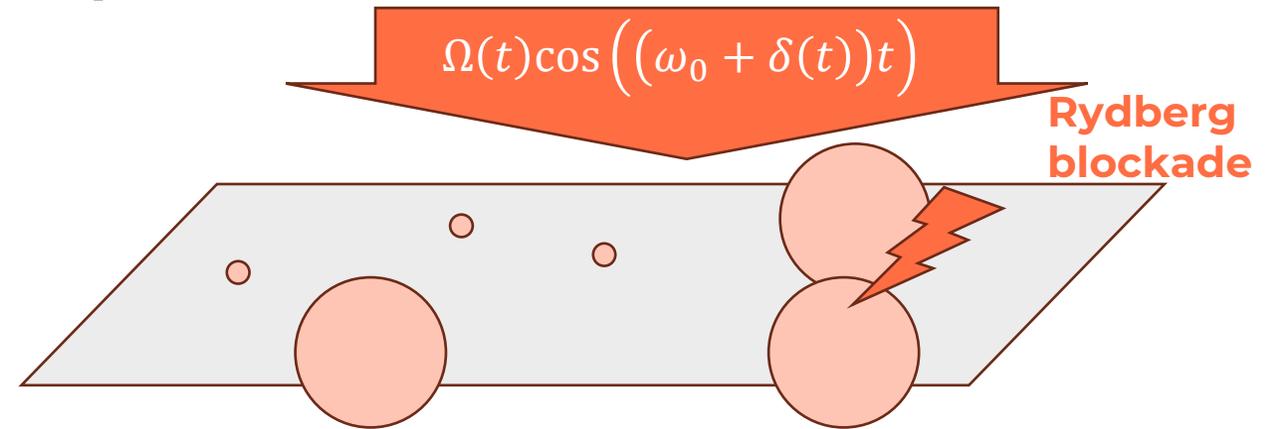
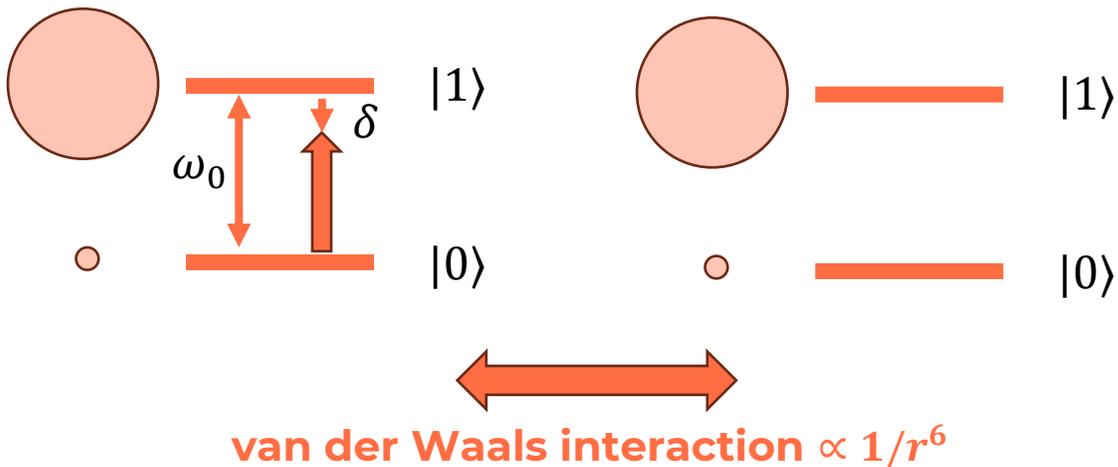
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detuning

vdW

Rabi

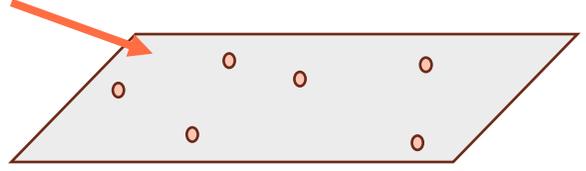
Differences: interaction term

- cannot turn vdW on/off
- $\frac{1}{r^6}$ dependence

Summary: analog quantum computation with Rydberg atoms

$$H_{\text{Ryd}} = -\delta(t) \sum_i \hat{n}_i + \sum_{i,j} \frac{C}{(r_i - r_j)^6} \hat{n}_i \hat{n}_j + \frac{\Omega(t)}{2} \sum_i \hat{\sigma}_i^x$$

Unit-disk graph



Detuning $\delta(t)$



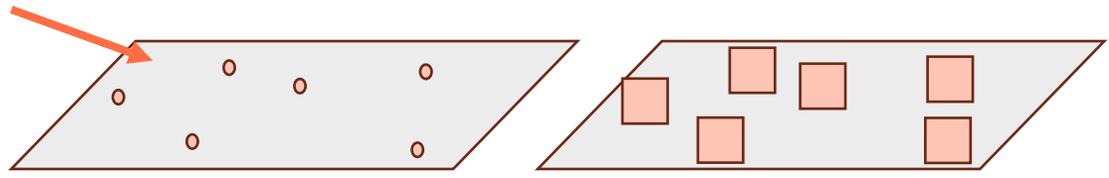
Rabi $\Omega(t)$



Summary: analog quantum computation with Rydberg atoms

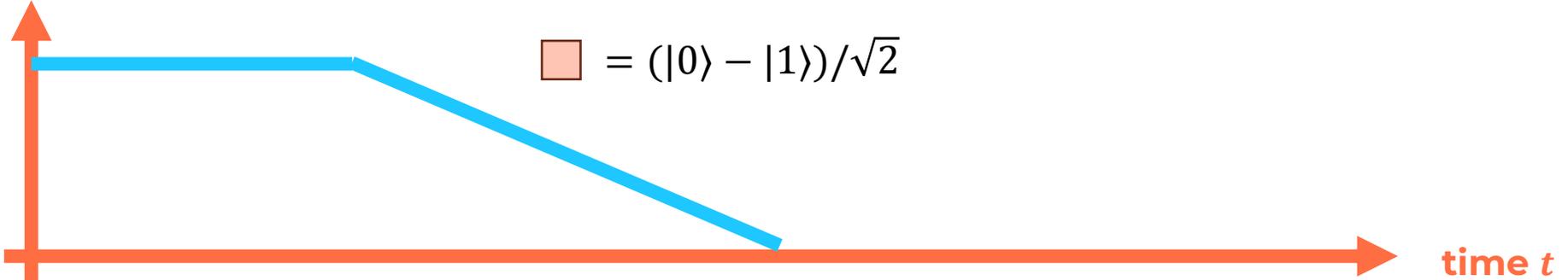
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Unit-disk graph

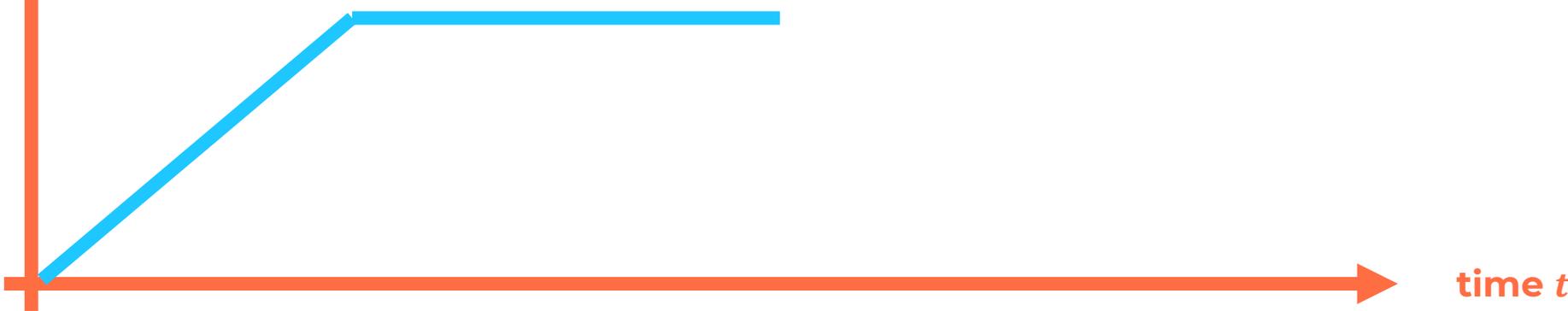


= $(|0\rangle - |1\rangle)/\sqrt{2}$

Detuning $\delta(t)$



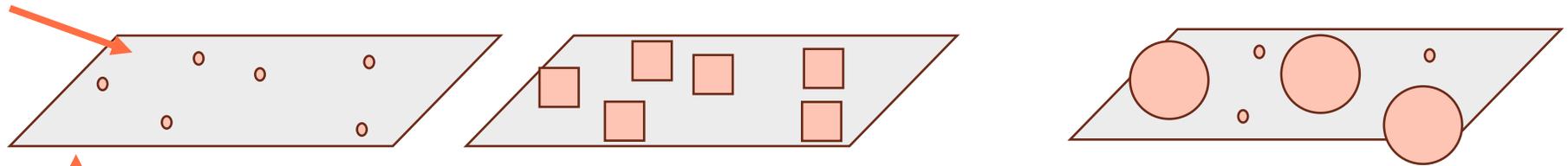
Rabi $\Omega(t)$



Summary: analog quantum computation with Rydberg atoms

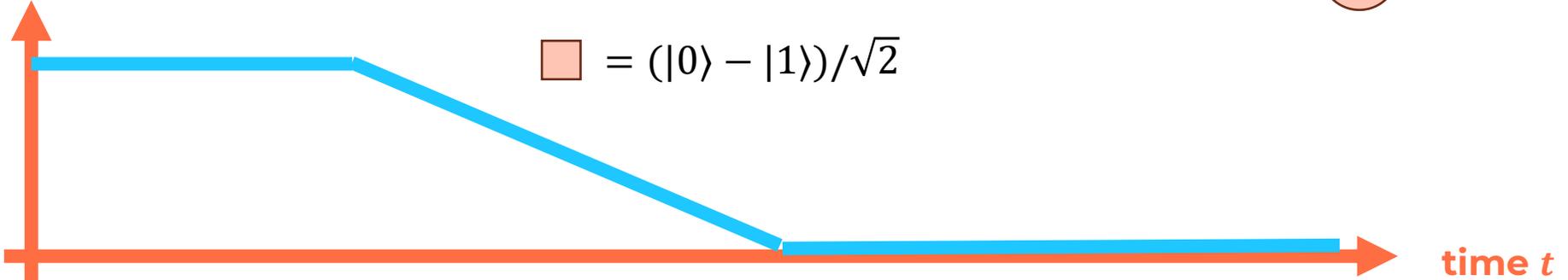
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Unit-disk graph

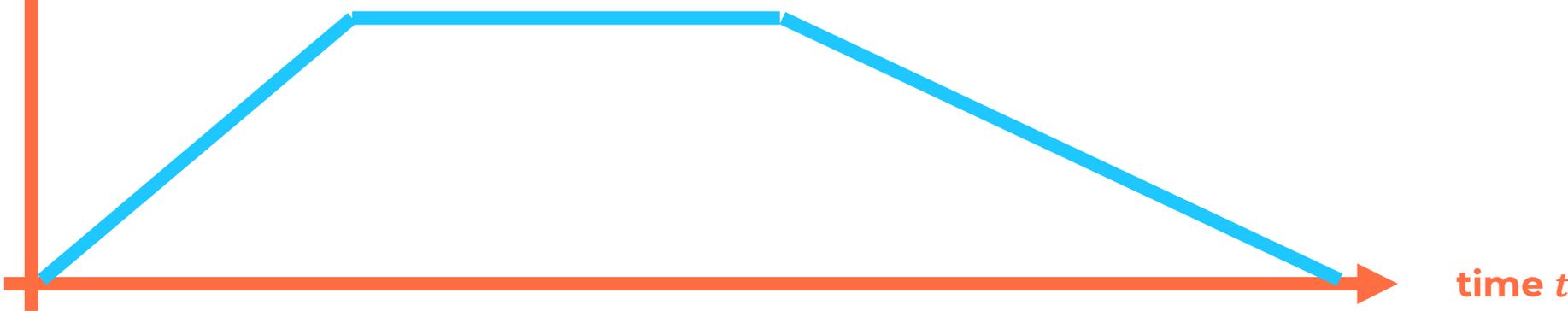


$$\square = (|0\rangle - |1\rangle) / \sqrt{2}$$

Detuning $\delta(t)$



Rabi $\Omega(t)$

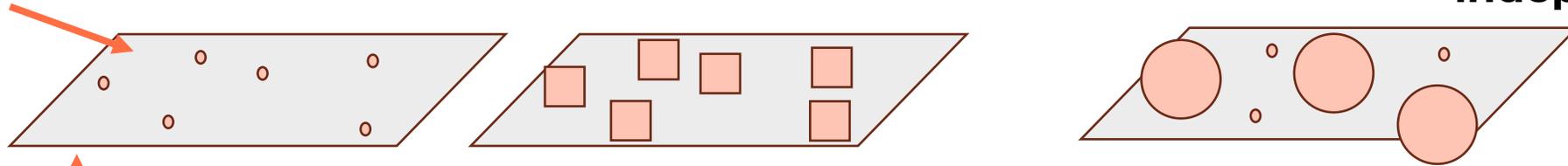


Summary: analog quantum computation with Rydberg atoms

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Unit-disk graph

Read off maximum independent set



$$\square = (|0\rangle - |1\rangle)/\sqrt{2}$$

Detuning $\delta(t)$

time t

Rabi $\Omega(t)$

time t

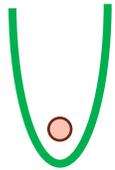
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2 The challenge of decoherence

Many things can go wrong!

$$H_{\text{Ryd}} = -\delta(t) \sum_i \hat{n}_i + \sum_{i,j} \frac{C}{(r_i - r_j)^6} \hat{n}_i \hat{n}_j + \frac{\Omega(t)}{2} \sum_i \hat{\sigma}_i^x$$

Optical tweezer



Heisenberg uncertainty:
 $r + \Delta r$

$$\Delta E_{\text{vdW}} \propto \frac{C}{(r_i - r_j)^7} (\Delta r_i + \Delta r_j)$$

Systematic (calibration) error: $\tilde{\delta}(t)$
+
Decoherence
 $\tilde{\delta}(t) + \Delta\delta(t)$

Systematic (calibration) error: $\tilde{\Omega}(t)$
+
Decoherence
 $\tilde{\Omega}(t) + \Delta\Omega(t)$

+ Readout errors

+ intrinsic limitations
cannot turn vdW on/off
 $\frac{1}{r^6}$ dependence
need long annealing time!

**Irreversible processes:
accumulate with time**

Can we make a quantitative prediction of the approximation ratio?

Numerical simulation of the analog computation

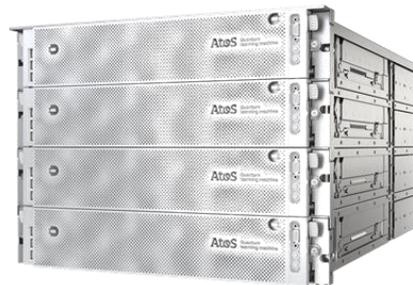
Simplified (but realistic) error model with

- **Readout noise.**
- **Dephasing white noise:** Lindblad equation

$$\frac{d\rho}{dt} = -i[H(t), \rho] - \frac{\gamma}{2} \sum_i \{n_i, \rho\} - 2n_i \rho n_i$$

- Solved with trajectories method (Dalibard, Castin & Mølmer 1992)

Eviden Qaptiva
compact
19" HPC appliance
NUMA architecture:
up to 32 Tb memory



Can we make a quantitative prediction of the approximation ratio?

Numerical simulation of the analog computation

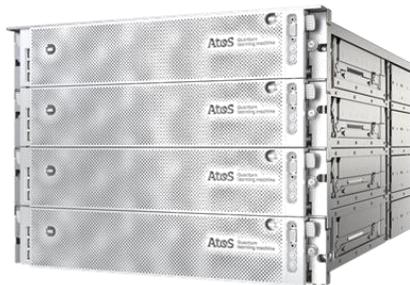
Simplified (but realistic) error model with

- **Readout noise.**
- **Dephasing white noise:** Lindblad equation

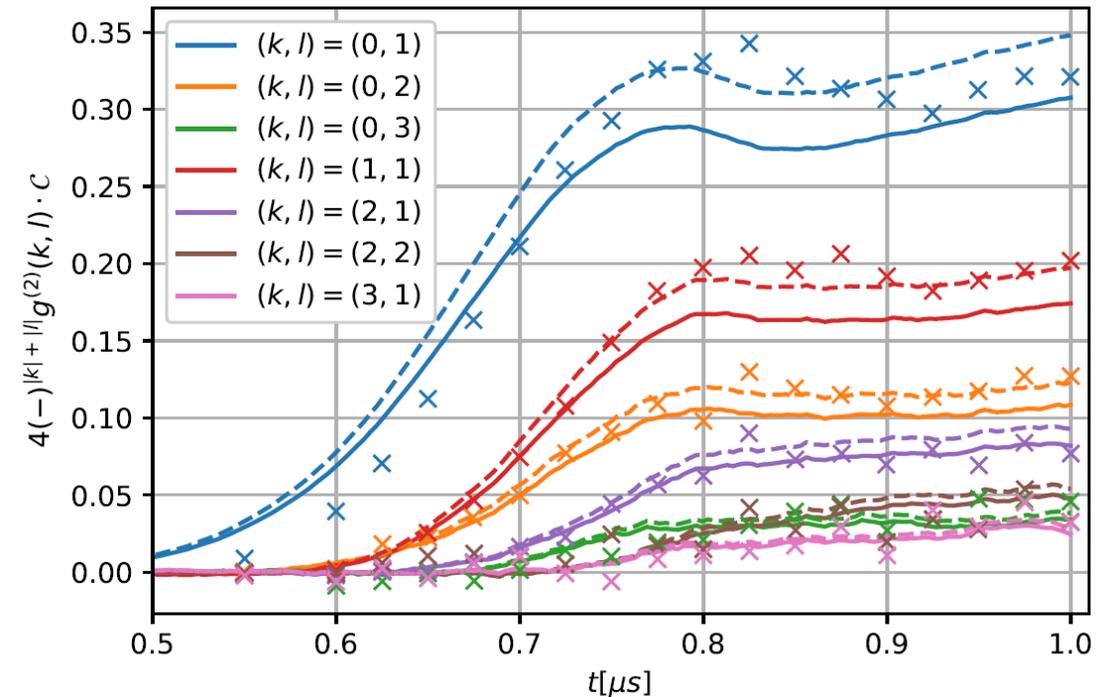
$$\frac{d\rho}{dt} = -i[H(t), \rho] - \frac{\gamma}{2} \sum_i \{n_i, \rho\} - 2n_i \rho n_i$$

- Solved with trajectories method (Dalibard, Castin & Mølmer 1992)

Eviden Qaptiva
compact
19" HPC appliance
NUMA architecture:
up to 32 Tb memory

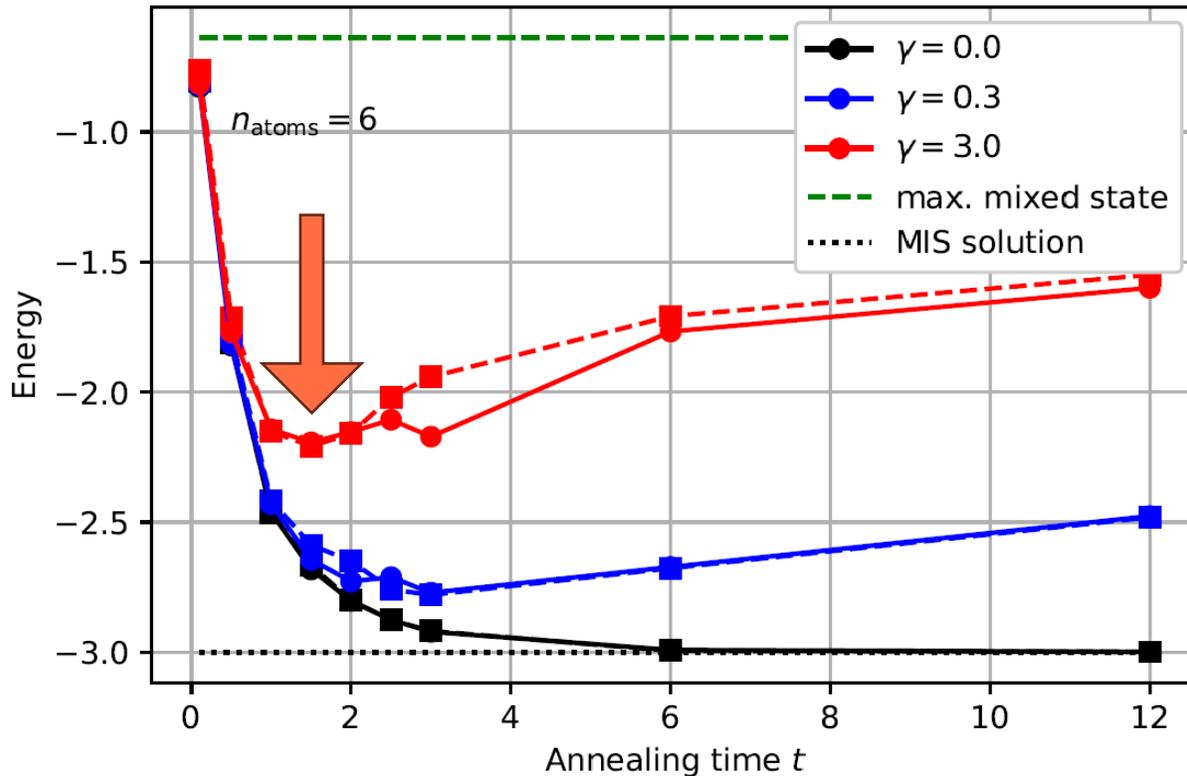


Validation for $\gamma = 3.0$ (exp: Lienhard '18)



Note: today's experiments: $\gamma = 0.3$ or even less!

Not too short, nor too long!

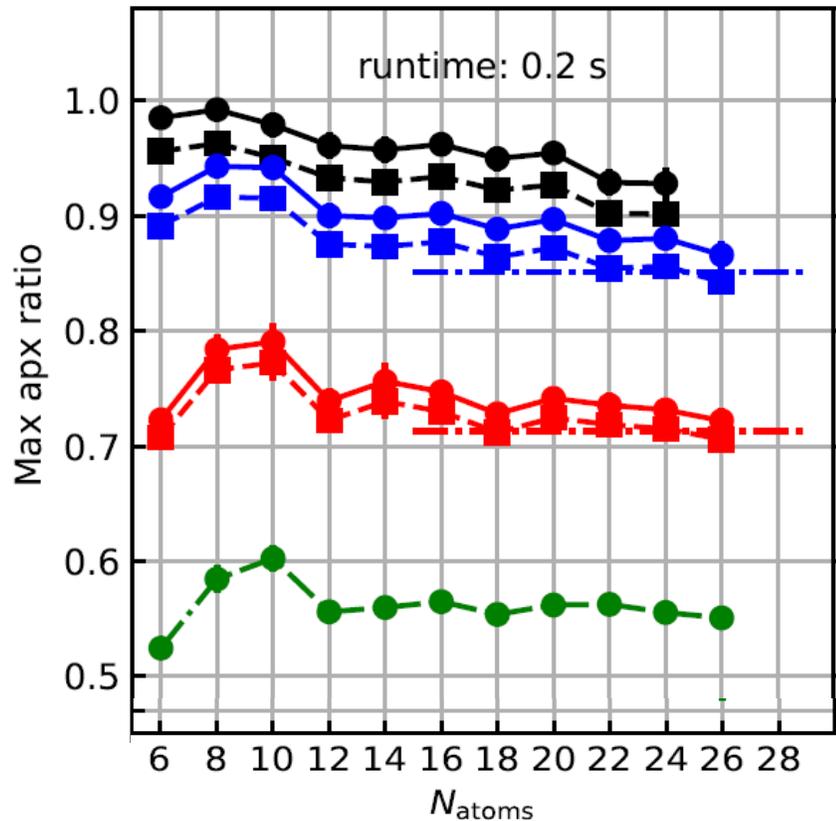
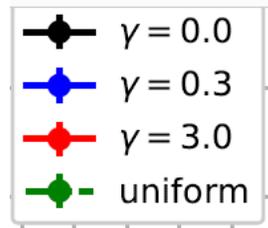


- **Noiseless case ($\gamma = 0$):** the longer, the better (adiabatic theorem)
- **Noisy case:** longer evolution = longer exposure to noise

In the following:

- work at optimal time
- work in “IS” subspace (dashed lines)

Dependence of expected apx ratio α_{avg} on noise γ and graph size (N_{atoms})?



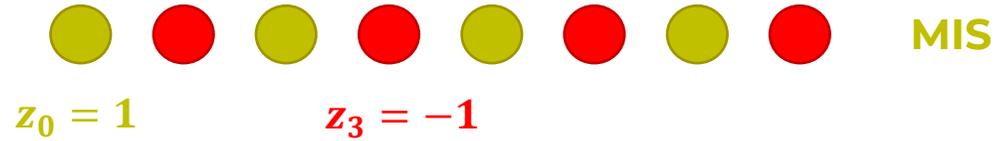
- **Noise γ :**

- Noise degrades α_{avg}
- Even at $\gamma = 0$, $\alpha_{avg} < 1$ (imperfect annealing schedule: VdW interactions...)

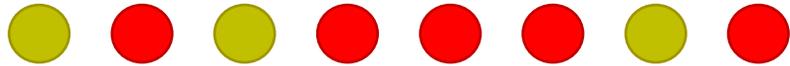
- **Size N_{atoms} :**

- Decreasing, then stable
- Stable earlier for higher γ

- ▶ Simple example: 1D chain



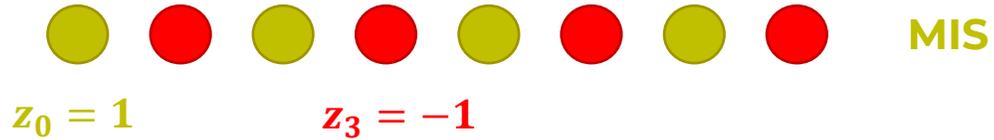
- ▶ Correlation function: $\langle z_i z_j \rangle$
 - Perfect “MIS” (antiferromagnetic) state: $\langle z_{2i} z_{2j} \rangle = 1, \langle z_{2i} z_{2j+1} \rangle = -1$
 - In the presence of decoherence: defects



Generically: $E(|\langle z_i z_{i+r} \rangle|) \propto e^{-r/\xi}$
with ξ correlation length.

Correlations & decoherence

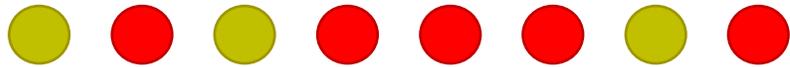
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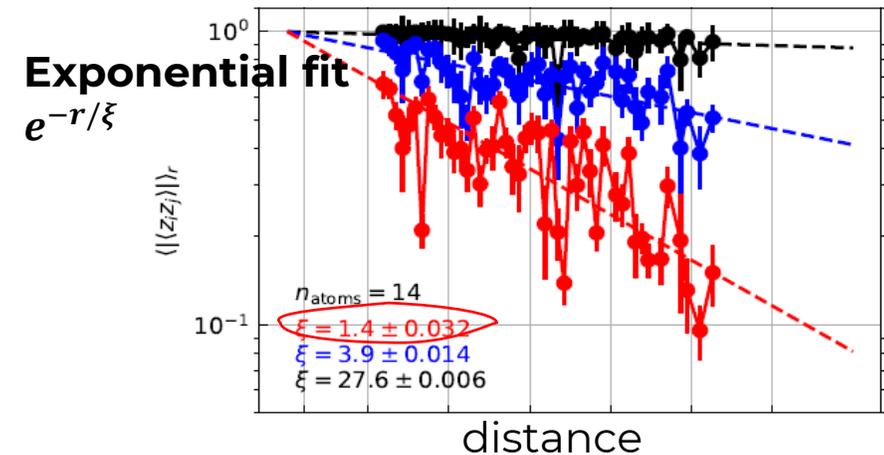
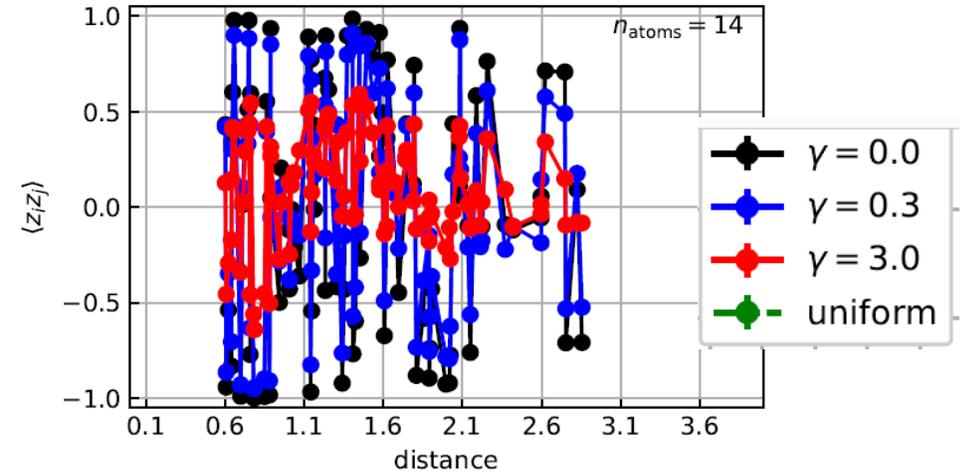
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Generically: $E(|\langle z_i z_{i+r} \rangle|) \propto e^{-r/\xi}$
with ξ correlation length.

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On our (2D) graphs:



- ▶ ξ decreases with increasing noise
- ▶ Can check: Correlation length is roughly independent of system size

Correlation length & approximation ratio

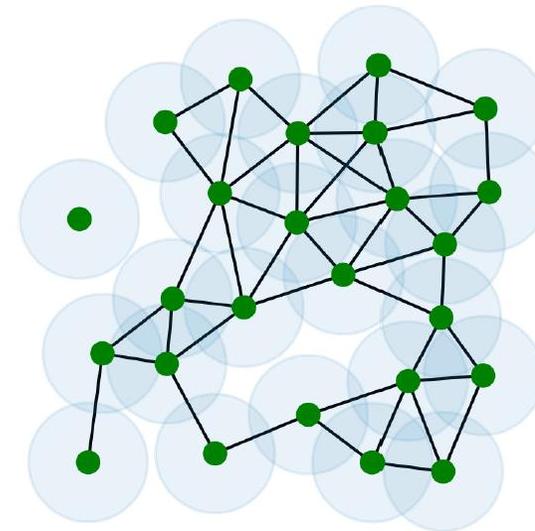
Can we relate the quality of a quantum algorithm with the correlation length?

Classical approaches... also have a kind of 'correlation length'!

All classical approaches:

“Divide and conquer”

Split the problem in smaller pieces and solve the pieces exactly



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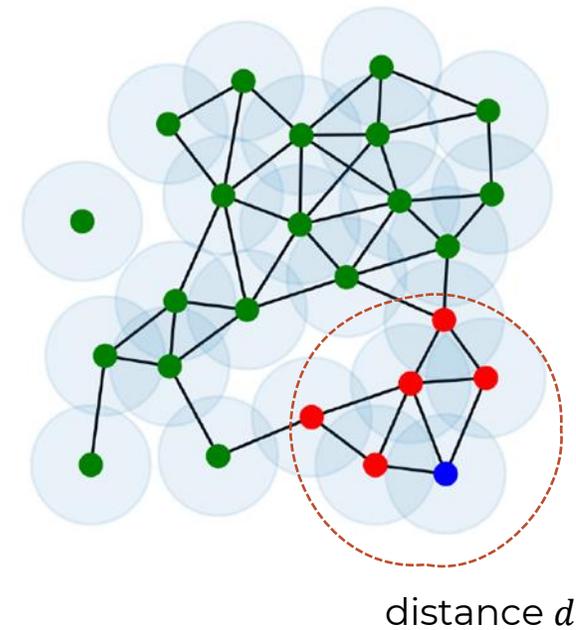
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For instance:

- Pick a **vertex** at random
- **Solve exactly** within 'distance' d



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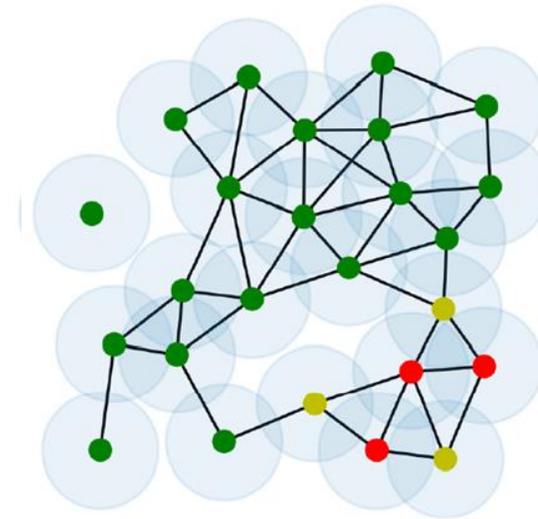
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- Remove subgraph (+connected vertices) from available vertices



MIS

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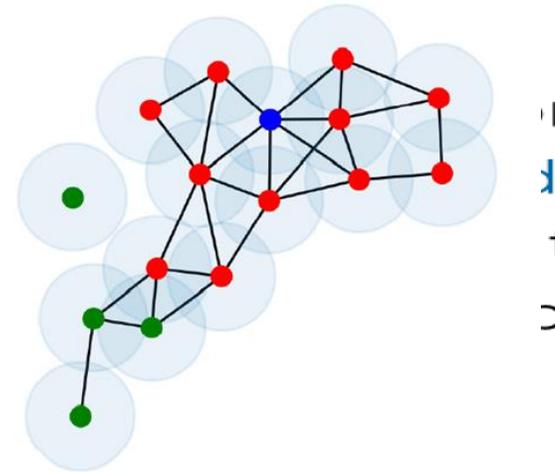
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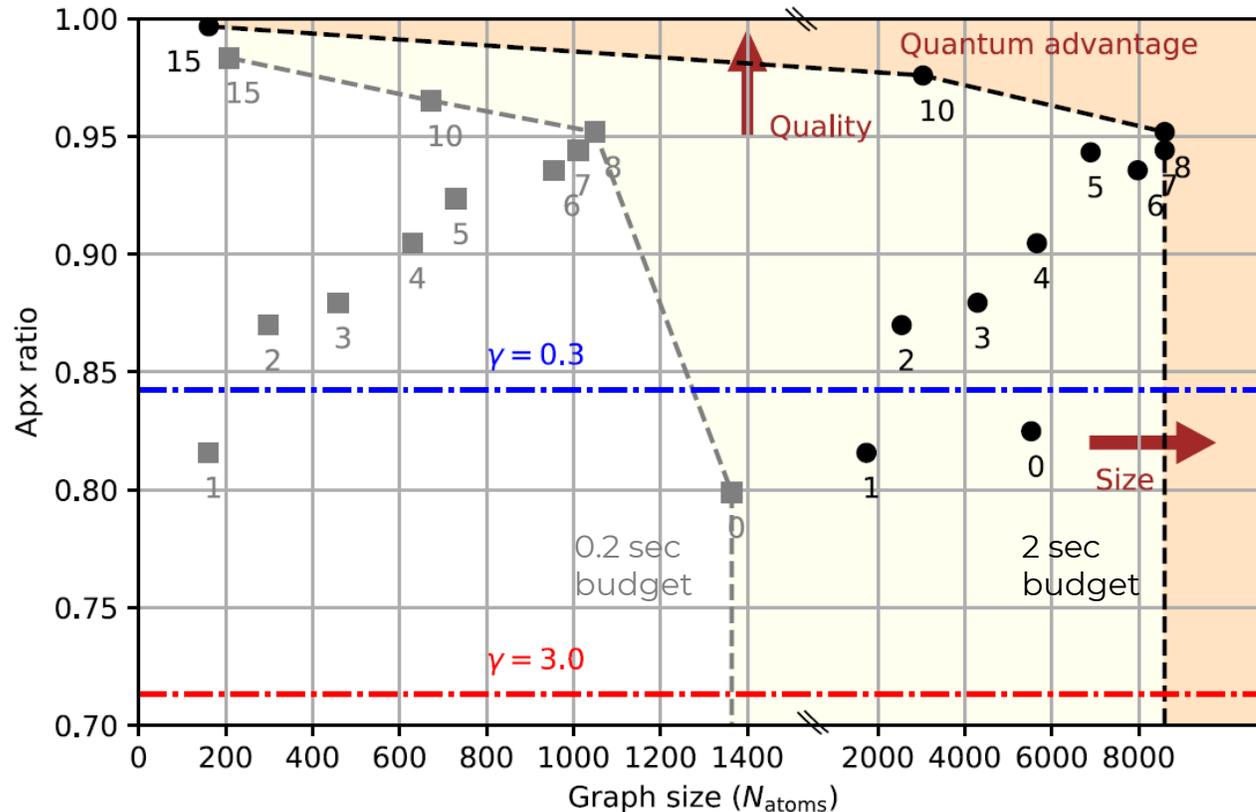
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- **Solve exactly** within 'distance' d
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- Remove subgraph (+connected vertices) from available vertices
- Iterate



A roadmap for quantum 'advantage'



Two ways to overperform the classical heuristic:

- **Go faster for bigger systems** (beat the exponential): "Size"
 ~ 1,000 atoms for 0.2 secs
 ~ 8,000 atoms for 2 secs
- **Reach higher approximation ratios:** "Quality"
 ~ e.g, 0.97 for 2000 atoms (2 secs)

Scaling of repetition rate?

Better hardware?

(better readout, lower noise, ...)

(Cf circular Rydberg atoms, [Nguyen et al 2018](#))

Better algorithms?

QAOA (digital), QA (better resource Hamiltonian), ...

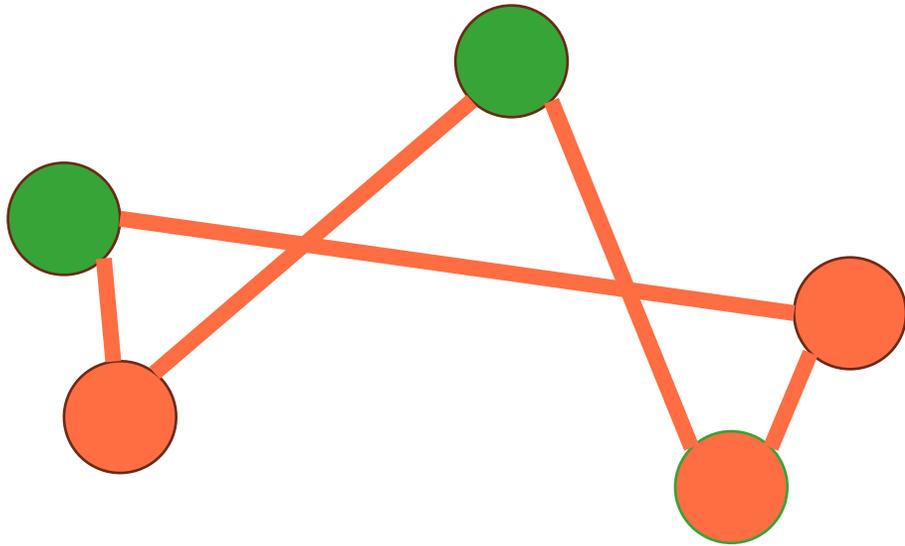
Or solve more difficult problems!

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3 Towards more general graphs... and applications

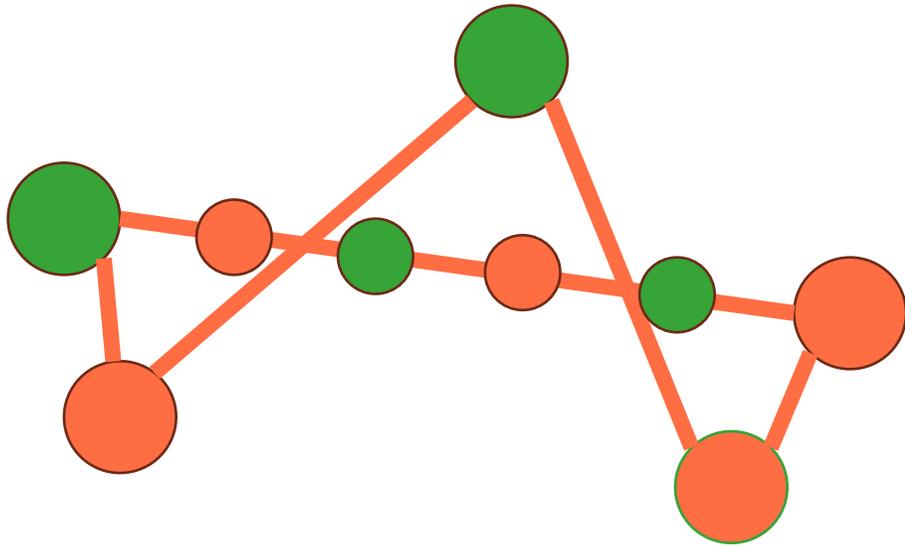
What about non-unit-disk graphs?

If G not unit disk graph: Harder problem... better candidate for quantum acceleration!?



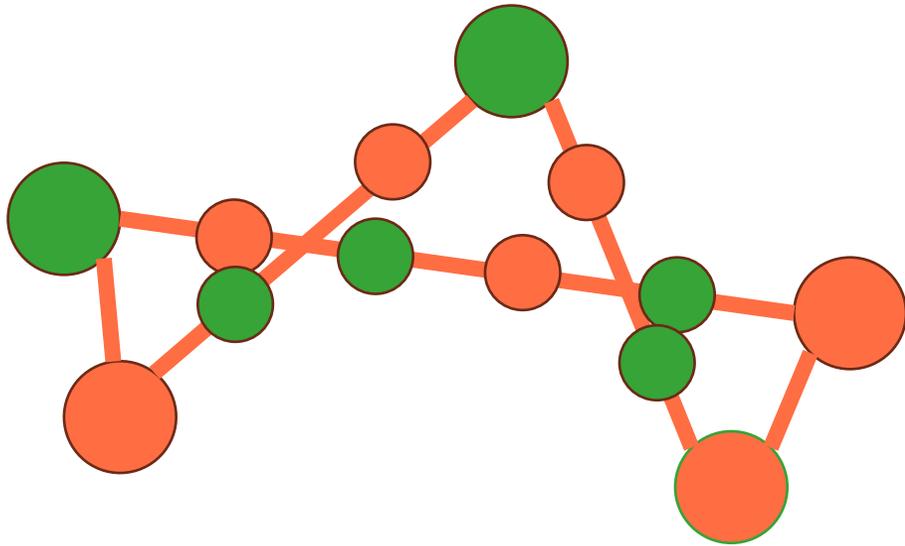
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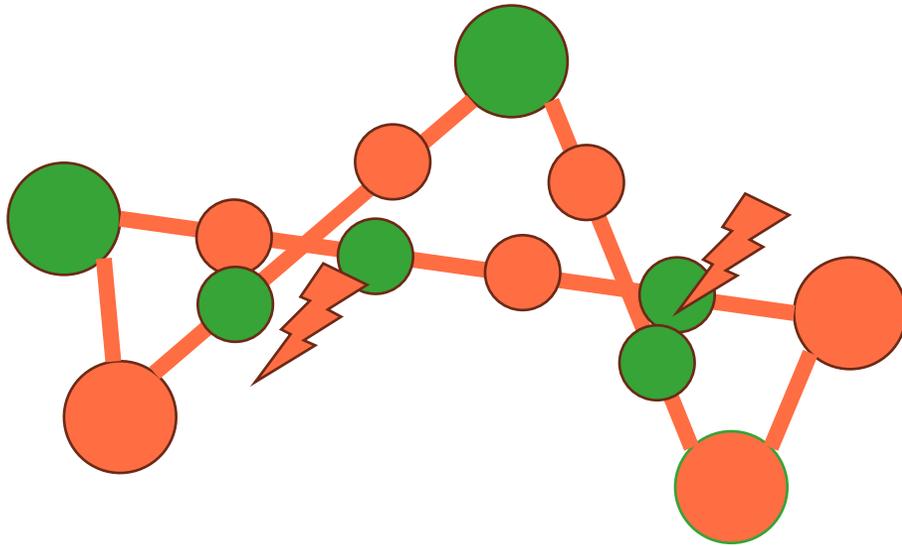
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Crossings of edges need to be handled properly!



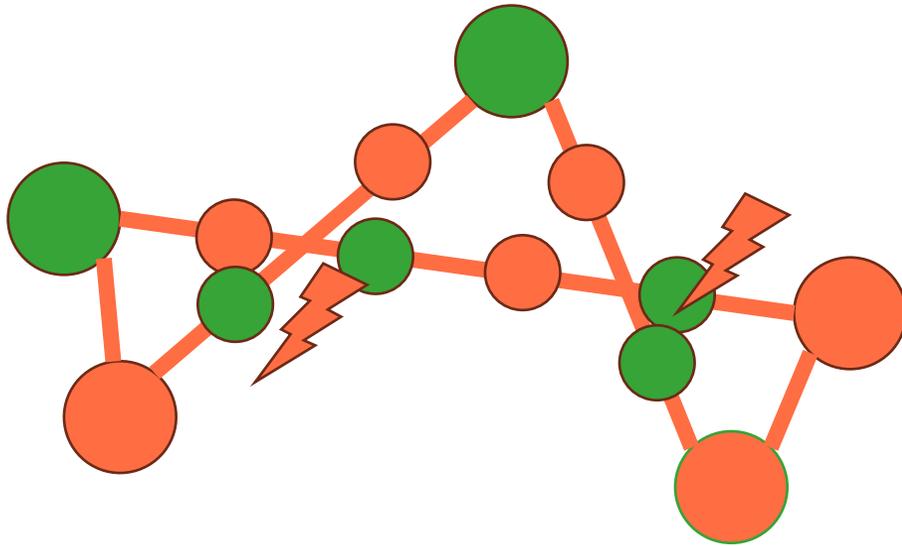
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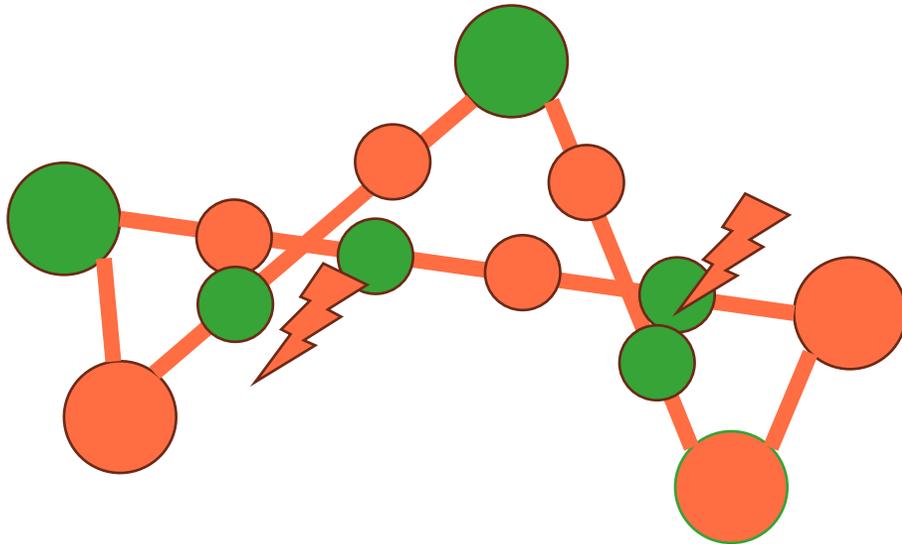
Nguyen et al 2022:

Can write a larger unit-disk graph G' such that:
the MIS of original graph G can be read off MIS of G' !



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Comes at a price:

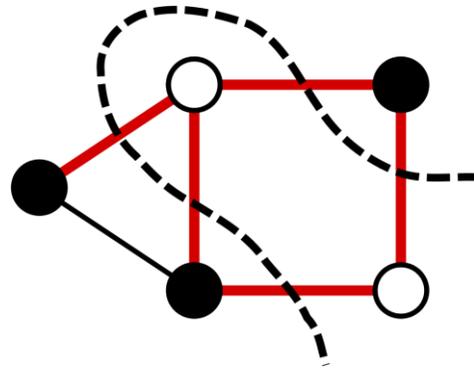
- Number of vertices (atoms): **quadratic increase: $O(N^2)$ atoms!**
- Needs local detuning $\delta_i(t)$!

What about other combinatorial optimization problems?

At the cost of adding vertices, can solve any « **quadratic, unconstrained optimization problem** » (QUBO):

$$H = \sum_{ij} J_{ij} s_i s_j + \sum_i h_i s_i \quad s_i \in \{-1, 1\}$$

For instance, **MaxCut problem**:



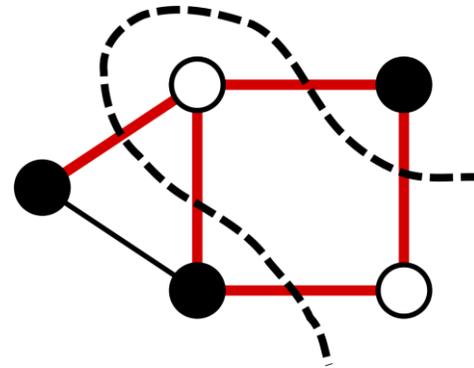
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Other example: **number factoring**.

Hard problems (spin glass physics)

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What about... solving fermionic problems?

Hubbard model: prototypical correlated electron problem.

- Quantum simulation by cold fermionic atoms!
- Rydberg atoms seem to be limited: described by spin Hamiltonian

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Can one not always turn fermions into spins?

Yes (Jordan-Wigner, etc...), but Hamiltonian is very different from Rydberg Hamiltonian!

$$c_i^\dagger c_j \rightarrow X_i Z_{i+1} \cdots Z_{j-1} X_j (+ \cdots)$$

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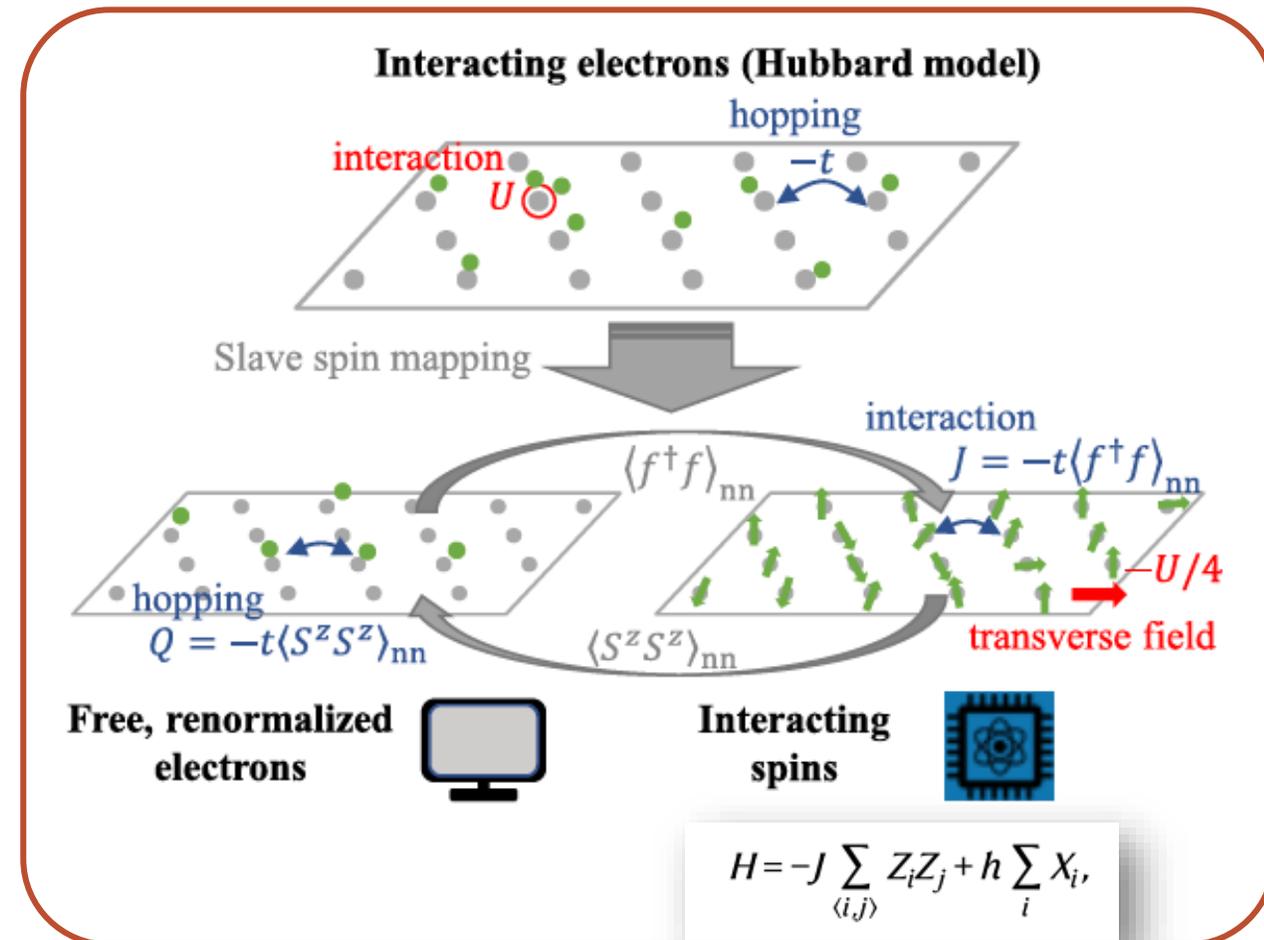
Yes (Jordan-Wigner, etc...), but Hamiltonian is very different from Rydberg Hamiltonian!

$$c_i^\dagger c_j \rightarrow X_i Z_{i+1} \cdots Z_{j-1} X_j (+ \dots)$$

Idea: use a slave-spin mapping: $c_{i\sigma}^\dagger = f_{i\sigma}^\dagger Z_i$

+ Mean-field decoupling

de' Medici 2005
Rüegg et al 2010
Hassan 2010



Very close to Rydberg Hamiltonian!

Effective model: Transverse Field Ising model (TFIM):

$$H_s^c = \sum_{i,j \in \mathcal{C}} J_{ij} S_i^z S_j^z + \frac{U}{4} \sum_{i \in \mathcal{C}} S_i^x + \sum_{i \in \mathcal{C}} h_i S_i^z,$$

... very close to Rydberg atom Hamiltonian!

$$\hat{H}_{\text{Rydberg}} = \sum_{i \neq j} \frac{C_6}{|\mathbf{r}_i - \mathbf{r}_j|^6} \hat{n}_i \hat{n}_j + \frac{\hbar \Omega(\tau)}{2} \sum_i \hat{S}_i^x - \hbar \delta(\tau) \sum_i \hat{n}_i,$$

Main challenges:

- Optimize atoms positions to reproduce J_{ij}
- Check robustness to decoherence

Using Rydberg atoms to deal with the spin model Michel, Henriet, Domain, Browaeys, TA, 2312.08065

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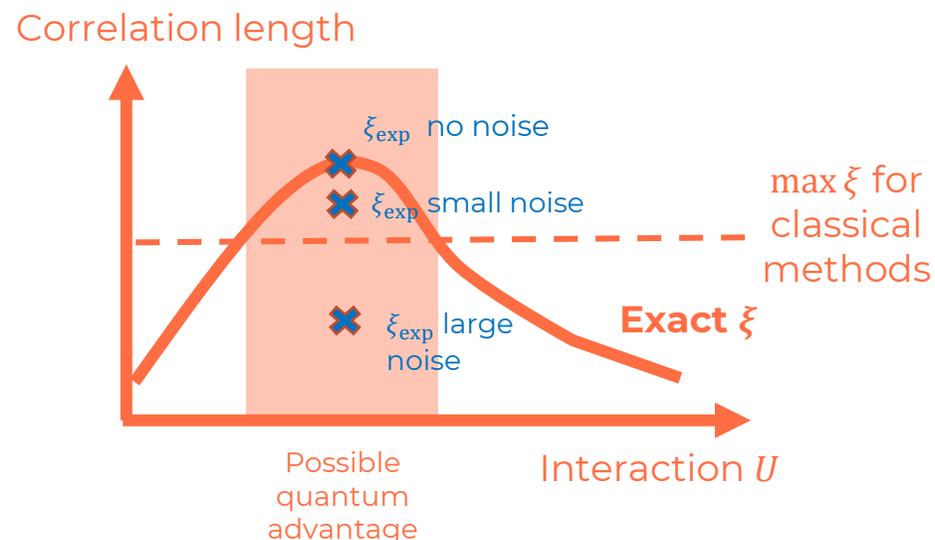
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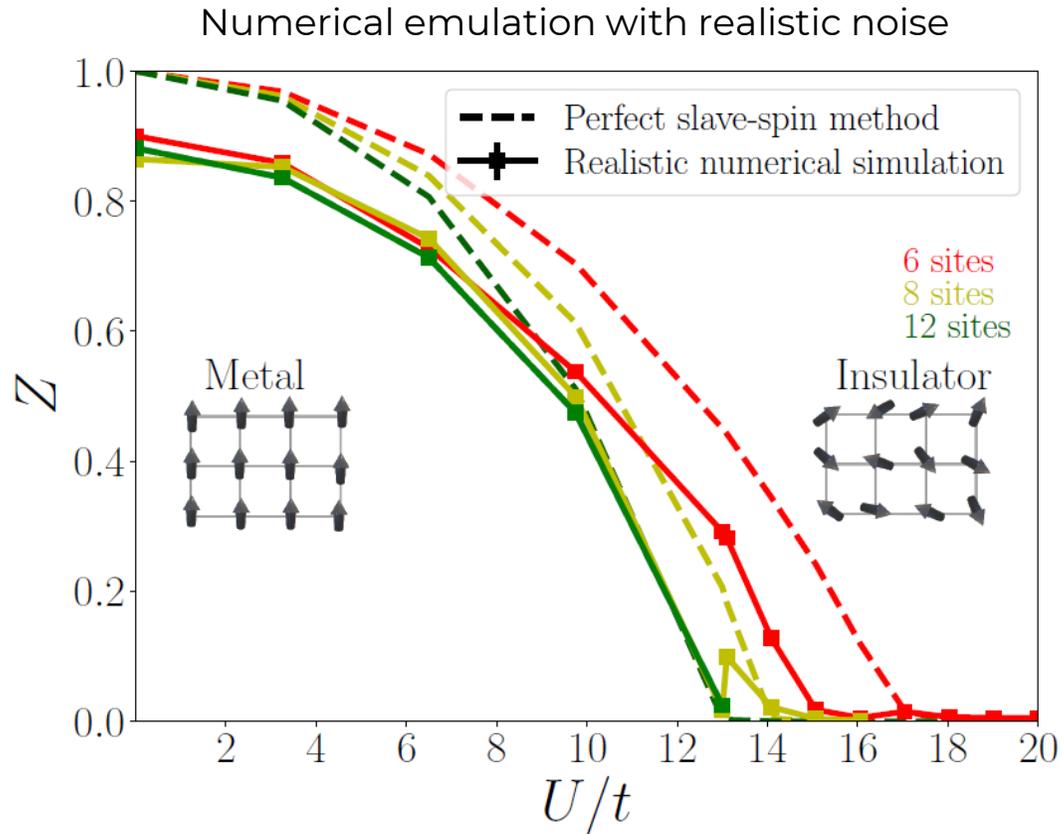
Hope for advantage w.r.t classical methods:



Mott physics with Rydberg atoms: (emulated) results

Can we locate Mott transition?

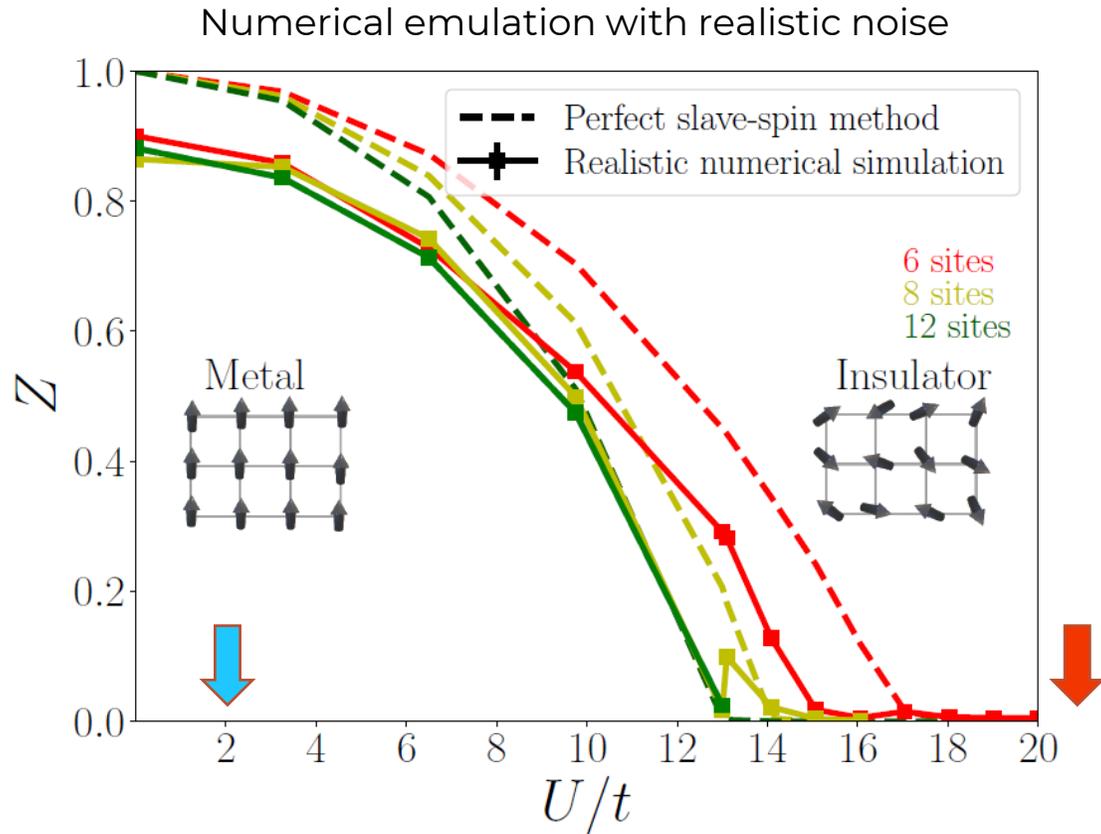
... in the presence of noise.



Mott physics with Rydberg atoms: (emulated) results

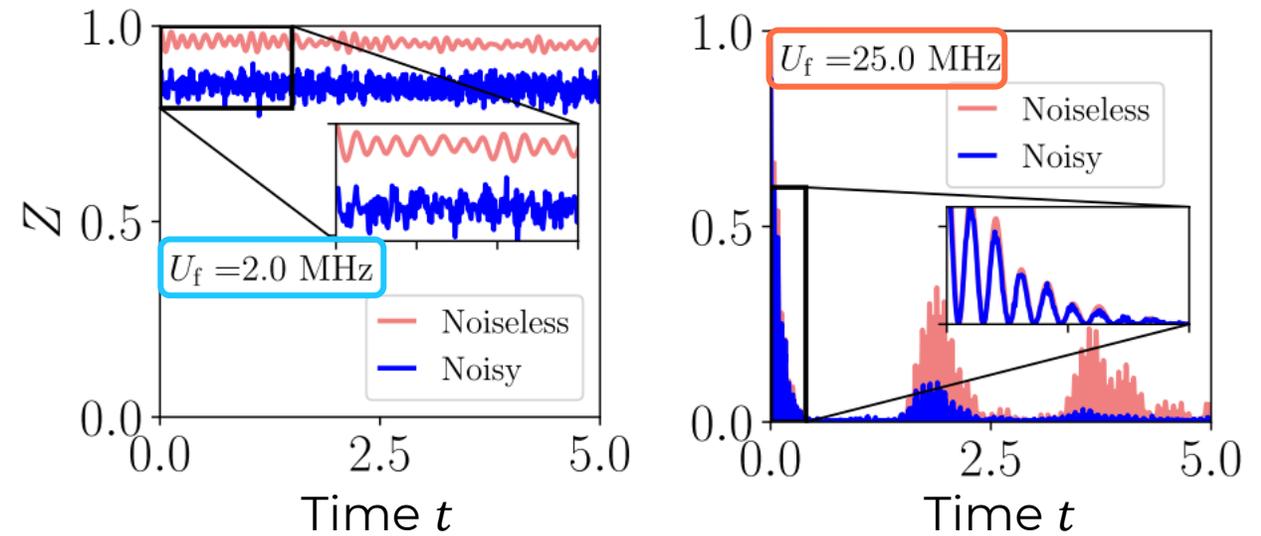
Can we locate Mott transition?

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What about out-of-equilibrium?

Interaction **quench** of the Hubbard model: becomes quench of transverse field in TFIM



Ongoing experimental implementation!

Conclusions

Rydberg platforms

- straightforward mapping to specific combinatorial optimization problems

Decoherence limits correlation length

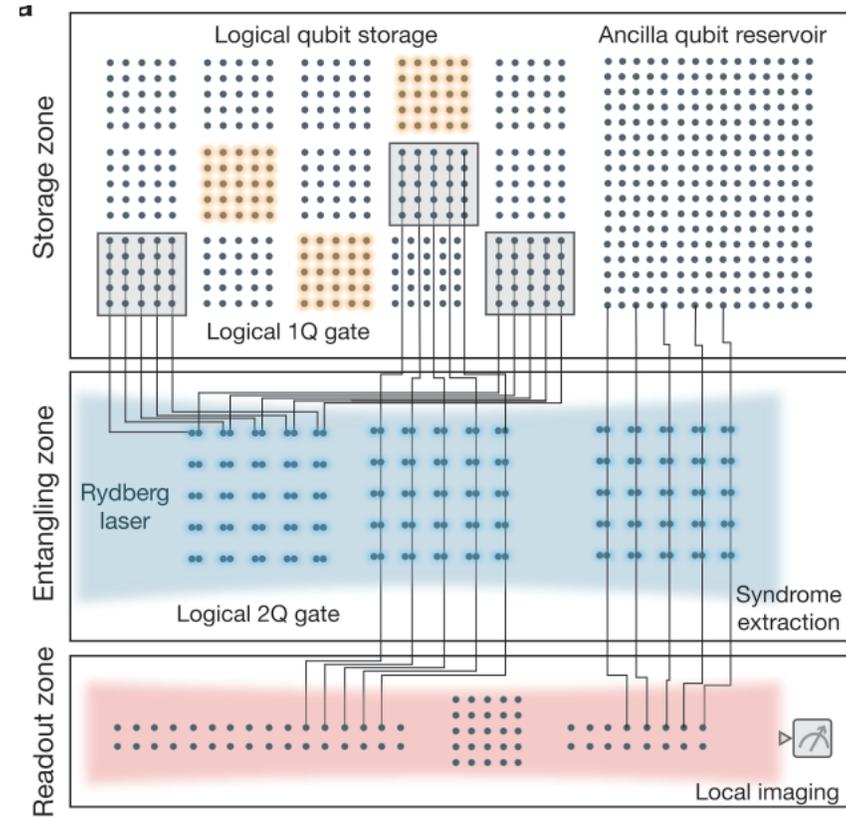
- Lesser success probability

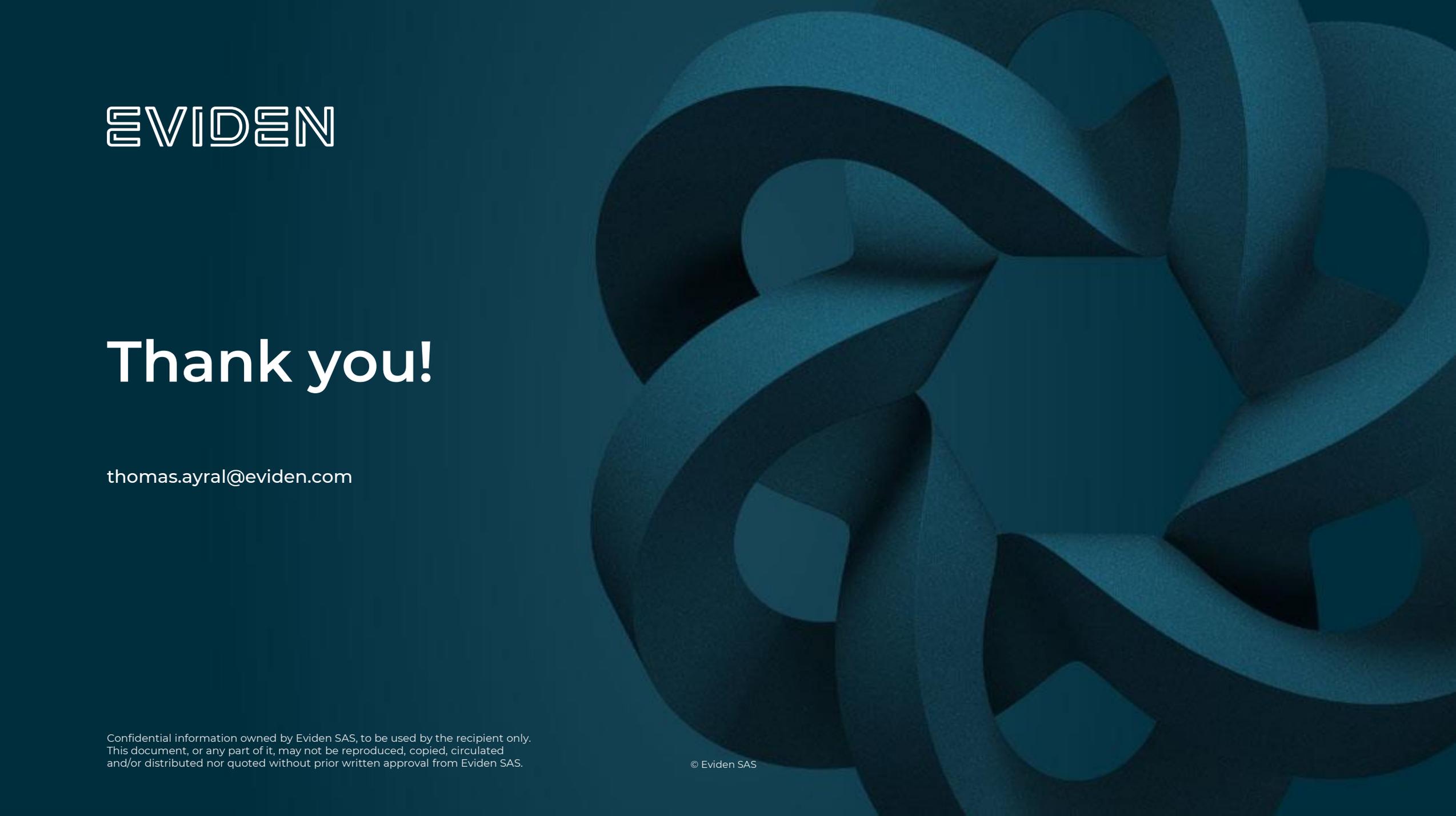
Recent extensions

- More general graphs
- Fermionic problems
- And others
- machine learning: quantum evolution kernel
[Henry et al 2021](#)

Recent breakthrough:

quantum error correction architecture ([Bluvstein et al 2024](#))





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Thank you!

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